Measures of Structural Complexity

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Measures of Complexity ... Measures of Structural Complexity:

Information Measures		Interpretation
Entropy Rate	h_{μ}	Intrinsic Randomness
Excess Entropy	\mathbf{E}	Info: Past to Future
Predictability Gain	G	Redundancy
Transient Information	\mathbf{T}	Synchronization

How related to statistical complexity C_{μ} ?

How to get from ϵM ?

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Measures from the ϵM :

Entropy Rate of a Process:

$$h_{\mu}(\Pr(\stackrel{\leftrightarrow}{S})) = \lim_{L \to \infty} \frac{H(L)}{L}$$

Directly from process's ϵM :

$$h_{\mu}\left(\Pr(\overset{\leftrightarrow}{S})\right) = h_{\mu}(\mathcal{S})$$

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Measures from the ϵM :

Entropy Rate given ϵM :

$$h_{\mu}(\mathcal{S}) = -\sum_{\mathcal{S} \in \mathcal{S}} \Pr(\mathcal{S}) \sum_{s \in \mathcal{A}, \mathcal{S}' \in \mathcal{S}} T_{\mathcal{S}\mathcal{S}'}^{(s)} \log_2 T_{\mathcal{S}\mathcal{S}'}^{(s)}$$

where $\Pr(S)$ is casual-state asymptotic probability.

Possible only due to ϵM unifilarity!

I-I mapping between measurement sequences & internal paths.

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Measures from the ϵM :

Entropy rate ...

Possible only due to unifilarity ...

What if you have a nonunifilar HMM for a process?

Consider two-state HMM presentation of SNS process:

$$1|\frac{1}{2} \underbrace{A} \underbrace{B} 1|\frac{1}{2}$$

$$\mathcal{B} = \{0, 1\}$$

$$T^{(0)} = \begin{pmatrix} 0 & 0 \\ \frac{1}{2} & 0 \end{pmatrix} \quad T^{(1)} = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \\ 0 & \frac{1}{2} \end{pmatrix}$$

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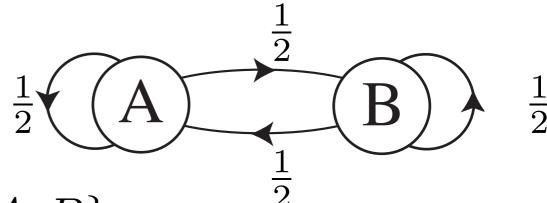
Monday, June 11, 2012 5

Measures from the ϵM :

Entropy rate ...

Possible only due to unifilarity ...

Internal Markov chain:



Internal (= Fair Coin): $A = \{A, B\}$

$$T = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \\ \frac{1}{2} & \frac{1}{2} \end{pmatrix} \quad \pi_V = \begin{pmatrix} \frac{1}{2}, \frac{1}{2} \end{pmatrix}$$

Entropy rate?

$$h_{\mu}(SNS) = -\sum_{v \in \mathcal{A}} \Pr(v) \sum_{v' \in \mathcal{A}} \Pr(v'|v) \log_2 \Pr(v'|v)$$

= 1

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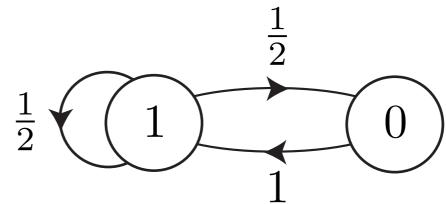
Monday, June 11, 2012 6

Measures from the ϵM :

Entropy rate ...

Possible only due to unifilarity ...

But support process is GMS



No consecutive 0s!

$$T = \begin{pmatrix} \frac{1}{2} & \frac{1}{2} \\ 1 & 0 \end{pmatrix} \quad \pi_V = (2/3, 1/3)$$

A restriction that lowers entropy rate.

$$h_{\mu}(GMS) = -\frac{2}{3} \sum_{s \in \mathcal{B}} \Pr(s|1) \log_2 \Pr(s|1)$$

= $\frac{2}{3}$

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Measures from the ϵM :

Entropy rate ...

Possible only due to unifilarity ...

$$h_{\mu}(SNS) \gg h_{\mu}(GMS)$$

"SNS entropy rate" larger than support process? No!

The 2-state presentation overestimates entropy rate: Internal process more random than observed process.

Lesson: Cannot use nonunifilar HMM representation of process to calculate entropy rate.

So how to compute SNS process's entropy rate?

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Measures from the ϵM :

Entropy rate ...

Lesson: Need ϵM to calculate entropy rate.

SNS example: Nontrivial, countably infinite ϵM .

$$h_{\mu} \approx 0.6778 \text{ bits/symbol}$$

Curious:

Even to estimate a process's intrinsic randomness, need to infer its structure.

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Measures from the $\epsilon M...$

Statistical Complexity of ϵM :

$$C_{\mu}(\mathcal{S}) = -\sum_{\mathcal{S} \in \mathcal{S}} \Pr(\mathcal{S}) \log_2 \Pr(\mathcal{S})$$

where Pr(S) is casual-state asymptotic probability.

Meaning:

Shannon information in the causal states.

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Measures from the $\epsilon M...$

Statistical Complexity of a Process:

$$C_{\mu}\left(\Pr(\stackrel{\leftrightarrow}{S})\right) = C_{\mu}(\mathcal{S})$$

Meaning:

The amount of historical information a process stores.

The amount of structure in a process.

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Measures from the $\epsilon M...$

Excess Entropy: Three versions, all equivalent for ID processes

$$\mathbf{E} = \lim_{L \to \infty} [H(L) - h_{\mu}L]$$

$$\mathbf{E} = \sum_{L=1}^{\infty} [h_{\mu}(L) - h_{\mu}]$$

$$\mathbf{E} = I[S; S]$$

How to get, given ϵM ?

Special cases: When ϵM is IID, periodic, or spin chain.

General case: Need a new framework.

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Measures from the $\epsilon M...$

Excess Entropy ...

 ϵM is IID:

$$\mathbf{E} = 0 \qquad C_{\mu} = 0$$

 $\epsilon \mathbf{M}$ is Period P:

$$\mathbf{E} = \log_2 P = C_{\mu}$$

 ϵM is range-R spin chain:

$$\mathbf{E} = H(R) - Rh_{\mu}$$

$$\mathbf{E} = C_{\mu} - Rh_{\mu}$$

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Measures from the $\epsilon M...$

Excess Entropy ...

Typically for Markov Chains:

$$\mathbf{E} < C_{\mu}$$

What can be said in general?

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Measures from the ϵM ...

Bound on Excess Entropy:

$$\mathbf{E} \leq C_{\mu}$$

Proof sketch:

(I)
$$\mathbf{E} = I[\overrightarrow{S}; \overleftarrow{S}] = H[\overrightarrow{S}] - H[\overrightarrow{S} \mid \overleftarrow{S}]$$

(2) Causal States:
$$H[\overrightarrow{S} \mid \overrightarrow{S}] = H[\overrightarrow{S} \mid \mathcal{S}]$$

(3)
$$\mathbf{E} = H[\overrightarrow{S}] - H[\overrightarrow{S} | \mathcal{S}]$$

$$= I[\overrightarrow{S}; \mathcal{S}]$$

$$= H[\mathcal{S}] - H[\mathcal{S} | \overrightarrow{S}]$$

$$\leq H[\mathcal{S}] = C_{\mu}$$

 \bigcirc

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Measures from the ϵM ...

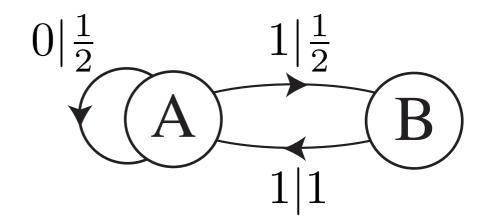
Bound on Excess Entropy ...

But, the bound is saturated!

Even process:

$$C_{\mu} = H(2/3) \approx 0.9182$$

$$\mathbf{E} \approx 0.9182$$



$$T^{(0)} = \begin{pmatrix} \frac{1}{2} & 0 \\ 0 & 0 \end{pmatrix}$$
$$T^{(1)} = \begin{pmatrix} 0 & \frac{1}{2} \\ 1 & 0 \end{pmatrix}$$
$$\pi_V = (2/3, 1/3)$$

When does this occur?

In general, need a new framework for answering this question.

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Measures from the ϵM ...

Bound on Excess Entropy ...

Consequence:

Can have $\mathbf{E} \to 0$ when $C_{\mu} \gg 1$.

(Cryptographic limit)

Excess entropy is not the process's stored information.

E is the apparent information, as revealed in measurement sequences.

Statistical complexity is stored information.

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Measures from the ϵM ...

Bound on Excess Entropy ...

Executive Summary:

 C_{μ} is the amount of information the process uses

to communicate

E bits of information from the past to the future.

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Measures from the ϵM ...

Bound on Excess Entropy: $\mathbf{E} \leq C_{\mu}$

Consequence:

The inequality is Why We Must Model.

Cannot simply use sequences as states.

There is internal structure not expressed by this.

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Process I-diagrams:

Process has an infinite number of RVs!

$$\Pr(\overrightarrow{X}) = \Pr(\dots, X_{-2}, X_{-1}, X_0, X_1, X_2, \dots)$$

Rather:
$$\Pr\left(\overleftrightarrow{X}\right) = \Pr\left(\overleftarrow{X} \ \overrightarrow{X}\right)$$

Start with 2-variable I-diagram and whittle down:

Past as composite random variable: X

Future as composite random variable: \overrightarrow{X}

Information measures:

$$H[\overleftarrow{X}] \ H[\overrightarrow{X}] \ H[\overrightarrow{X}, \overleftarrow{X}]$$

$$H[\overleftarrow{X}|\overrightarrow{X}] \ H[\overrightarrow{X}|\overleftarrow{X}] \ I[\overrightarrow{X}; \overleftarrow{X}] \ H[\overrightarrow{X}|\overleftarrow{X}] + H[\overleftarrow{X}|\overrightarrow{X}]$$

There are $3 = 2^2$ -I atomic information measures:

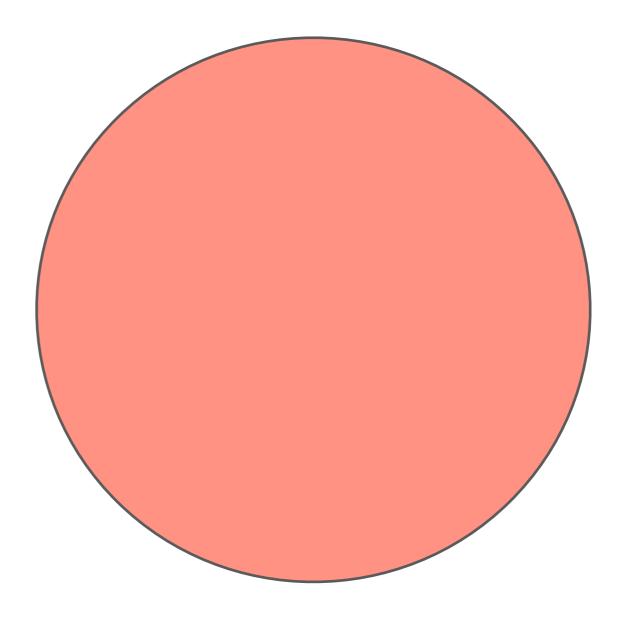
$$H[\overrightarrow{X}|\overrightarrow{X}] \quad H[\overleftarrow{X}|\overrightarrow{X}] \quad I[\overrightarrow{X};\overleftarrow{X}]$$

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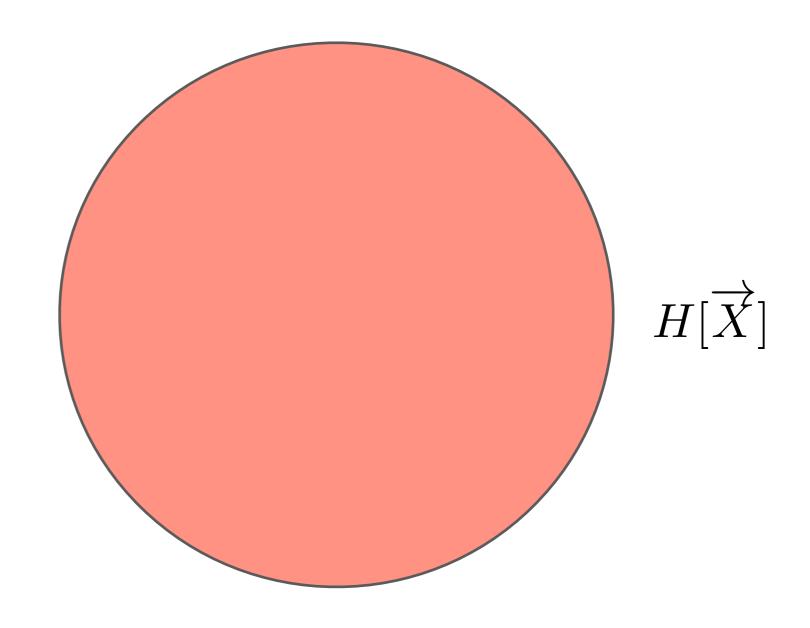
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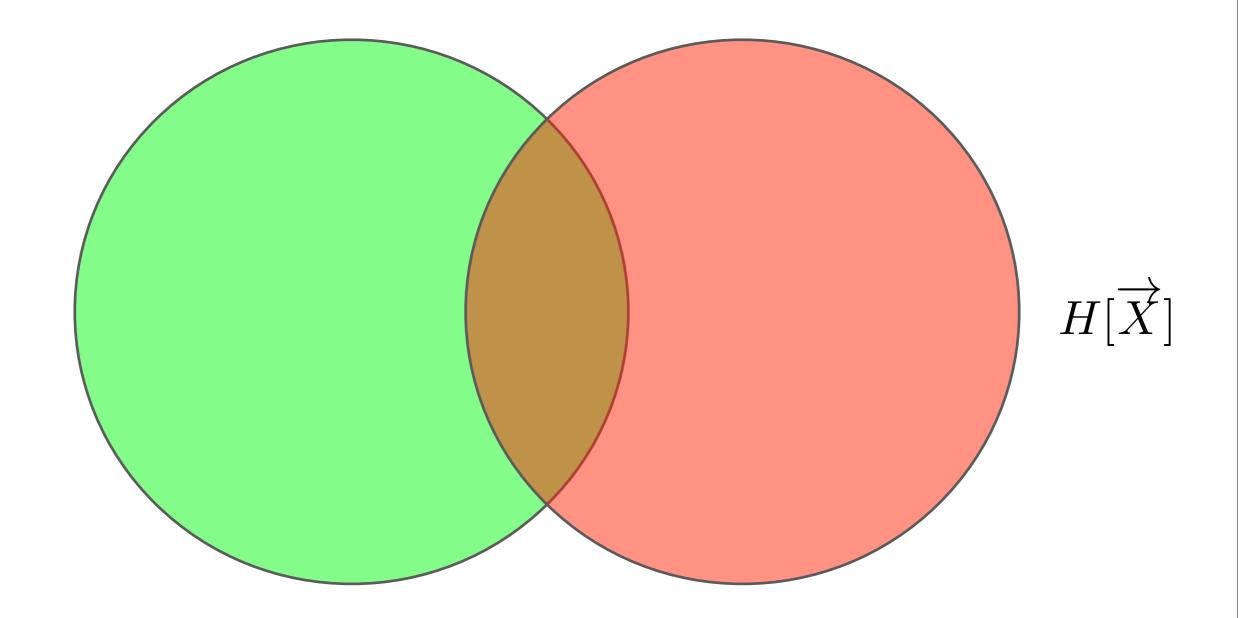


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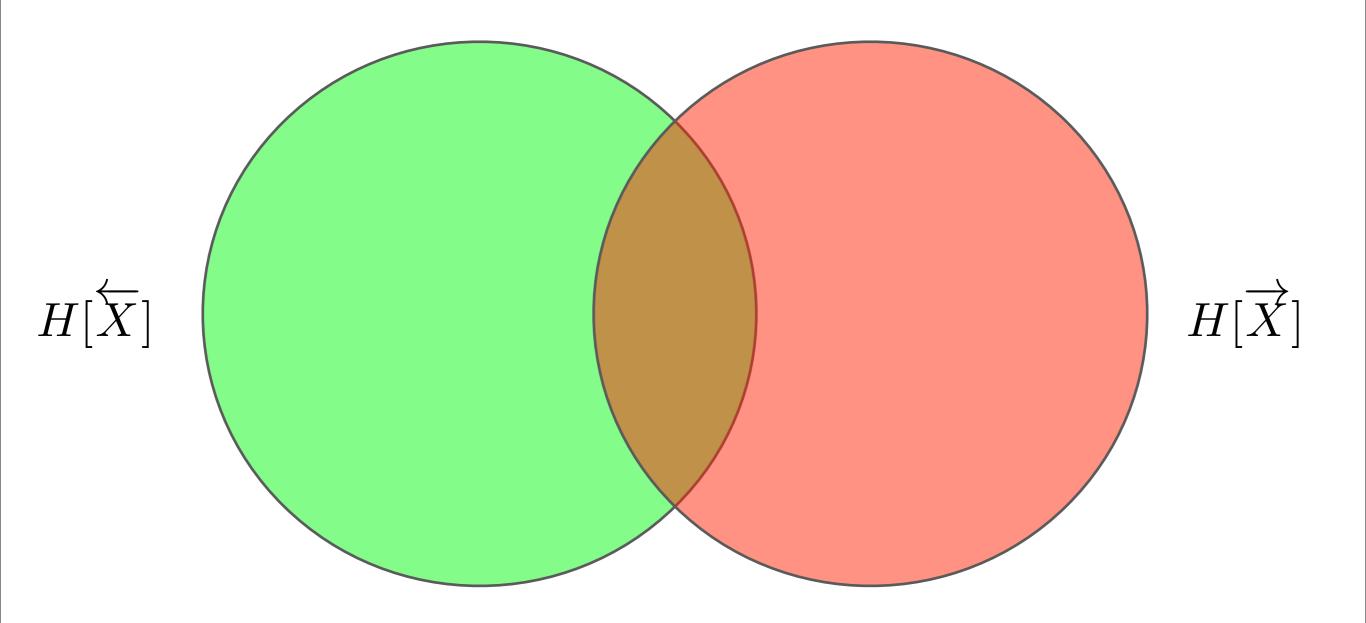
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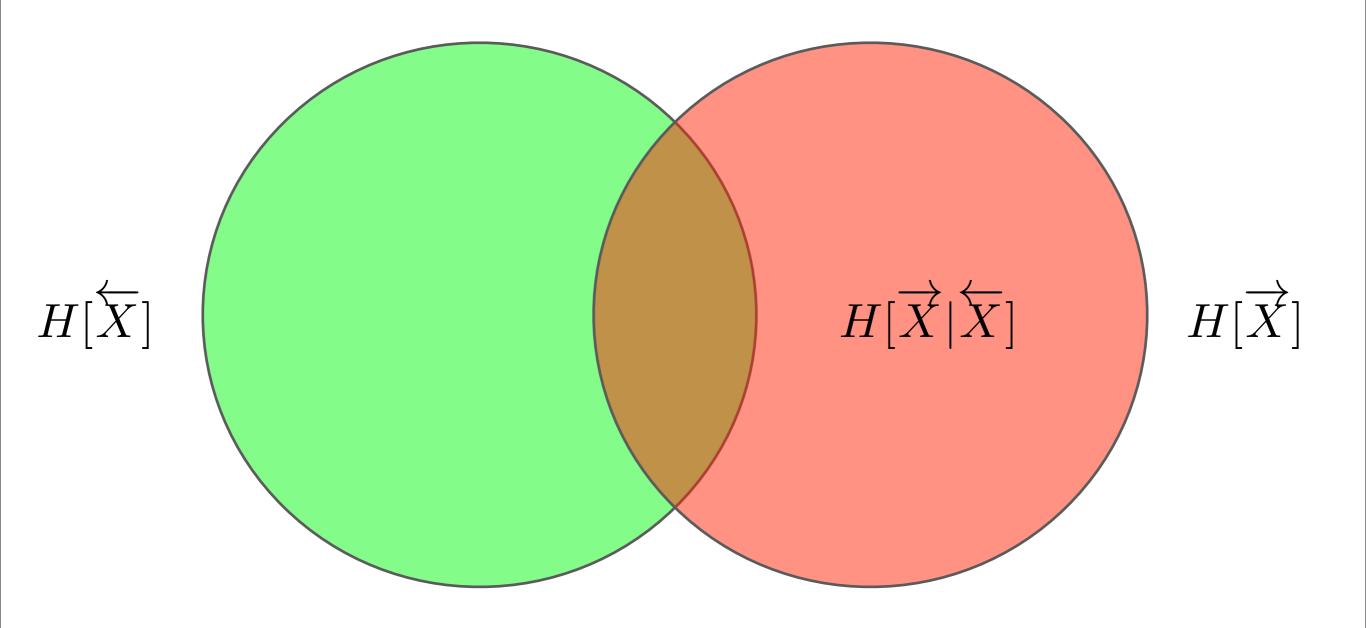
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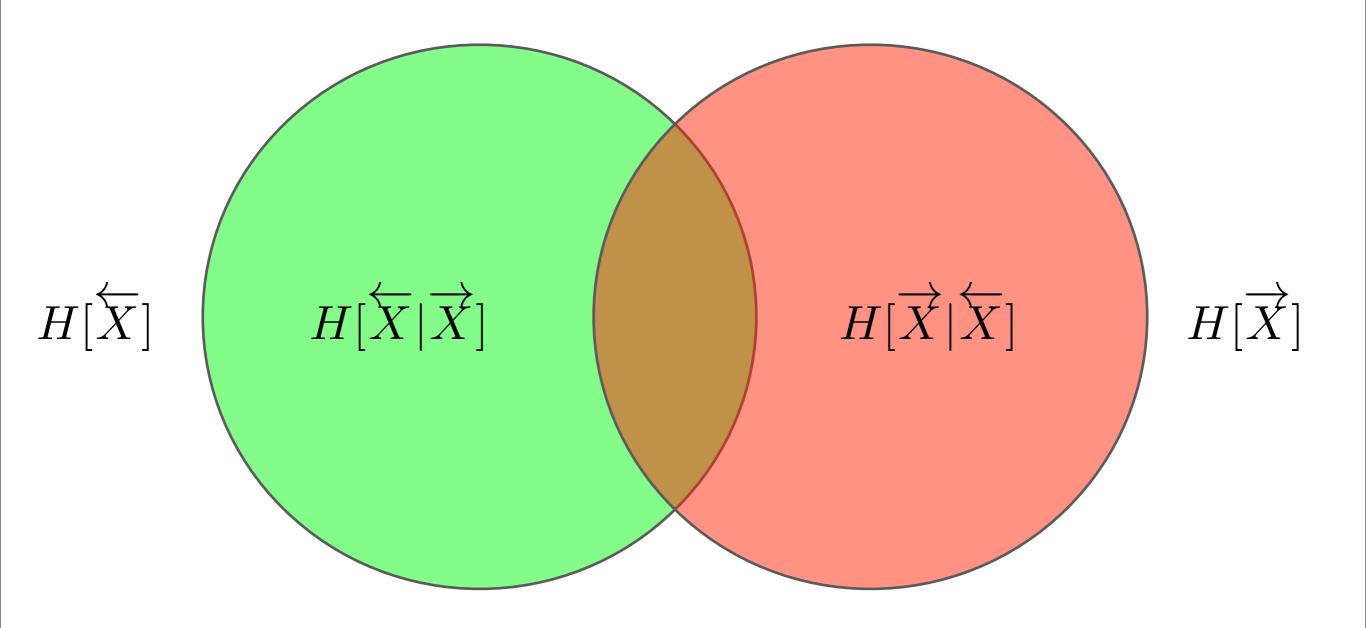
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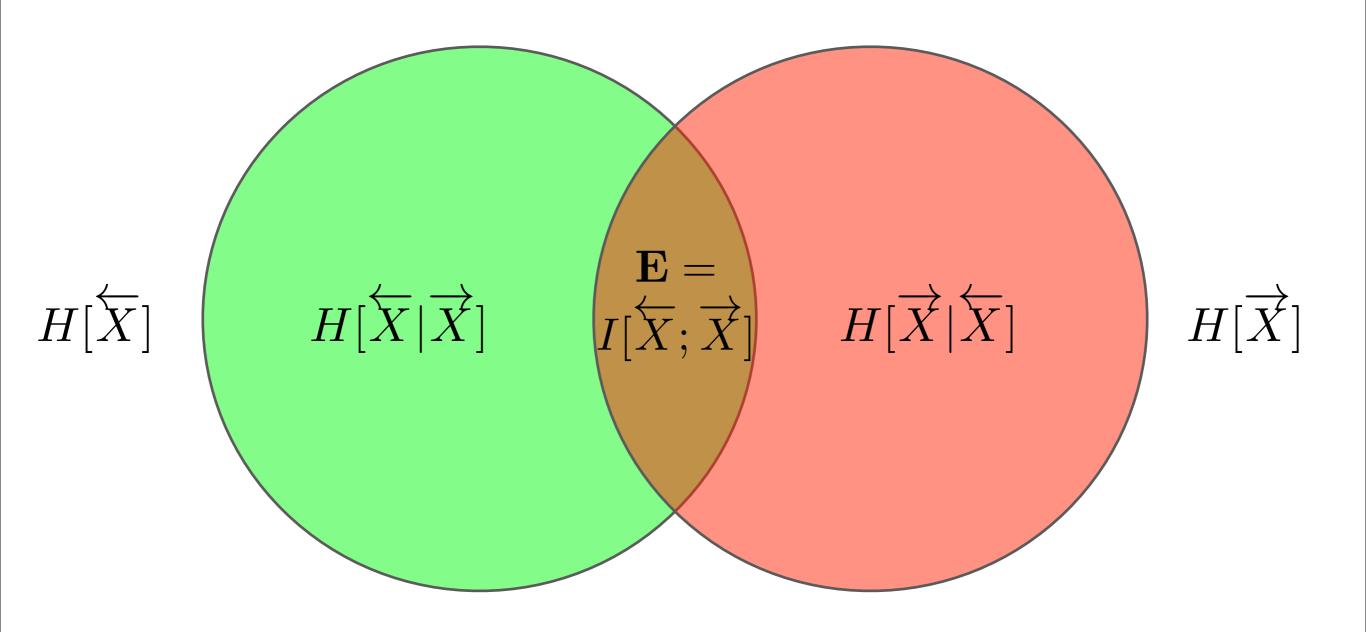
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Process I-diagram using E-machine:

Start with 3-variable I-diagram and whittle down:

Past as composite random variable: X

Future as composite random variable: \overrightarrow{X}

Causal states: $S \in S$

Information measures:

$$H[\overleftarrow{X}] \ H[\overrightarrow{X}] \ H[\mathcal{S}] \ \cdots \ I[\overrightarrow{X}; \overleftarrow{X}; \mathcal{S}] \ \cdots \ H[\overrightarrow{X}, \overleftarrow{X}, \mathcal{S}]$$

There are $8 = 2^3$ atomic information measures.

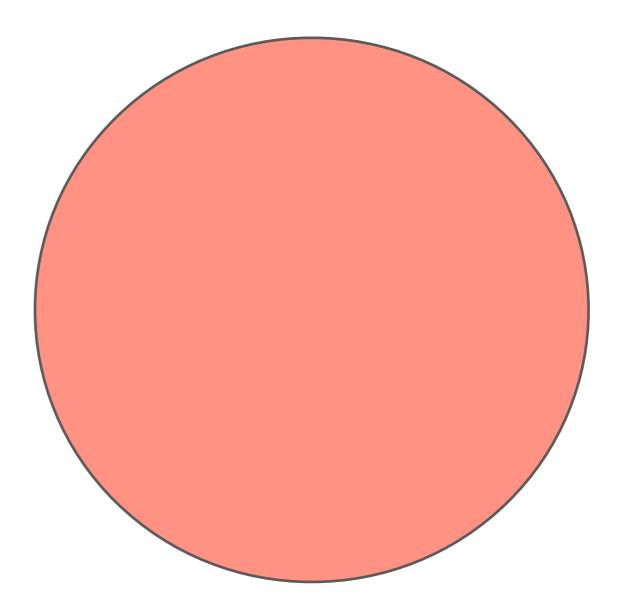
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Information Diagrams for Processes
Process I-diagram using \(\mathcal{E}\)-machine ...

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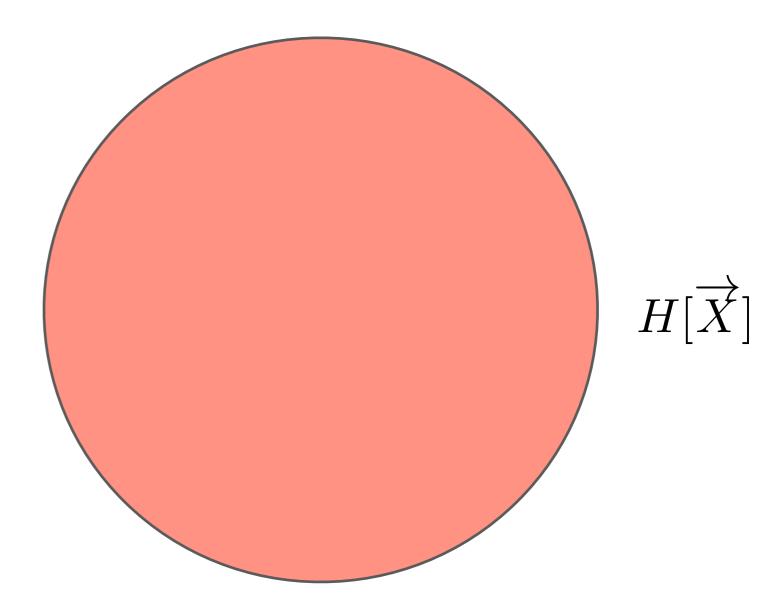
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Information Diagrams for Processes
Process I-diagram using \(\mathcal{E}\)-machine ...



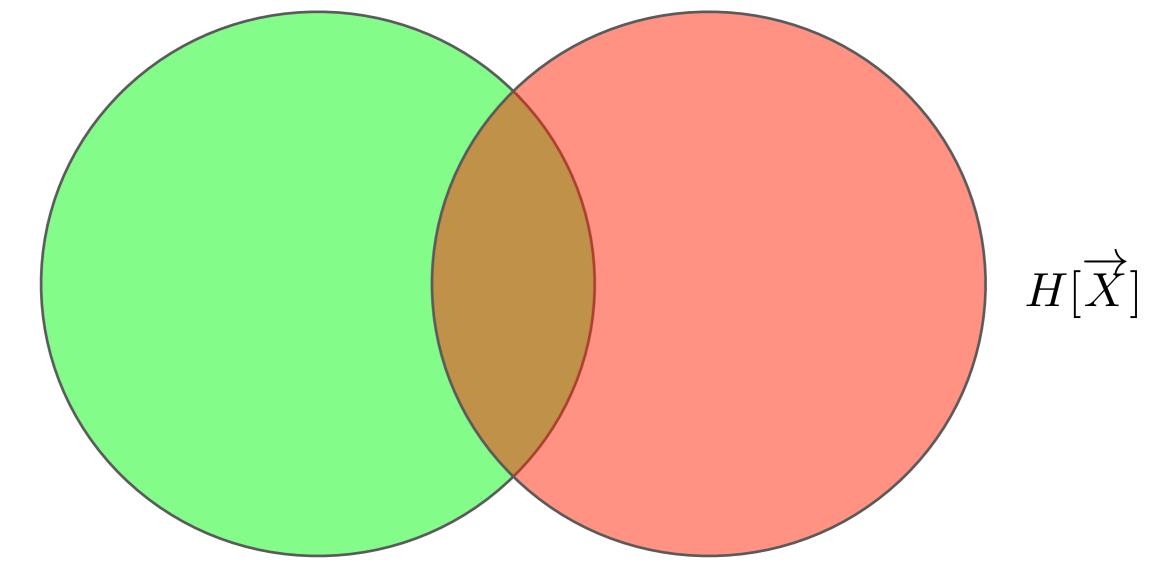
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Information Diagrams for Processes Process I-diagram using \(\mathcal{E}\)-machine ...



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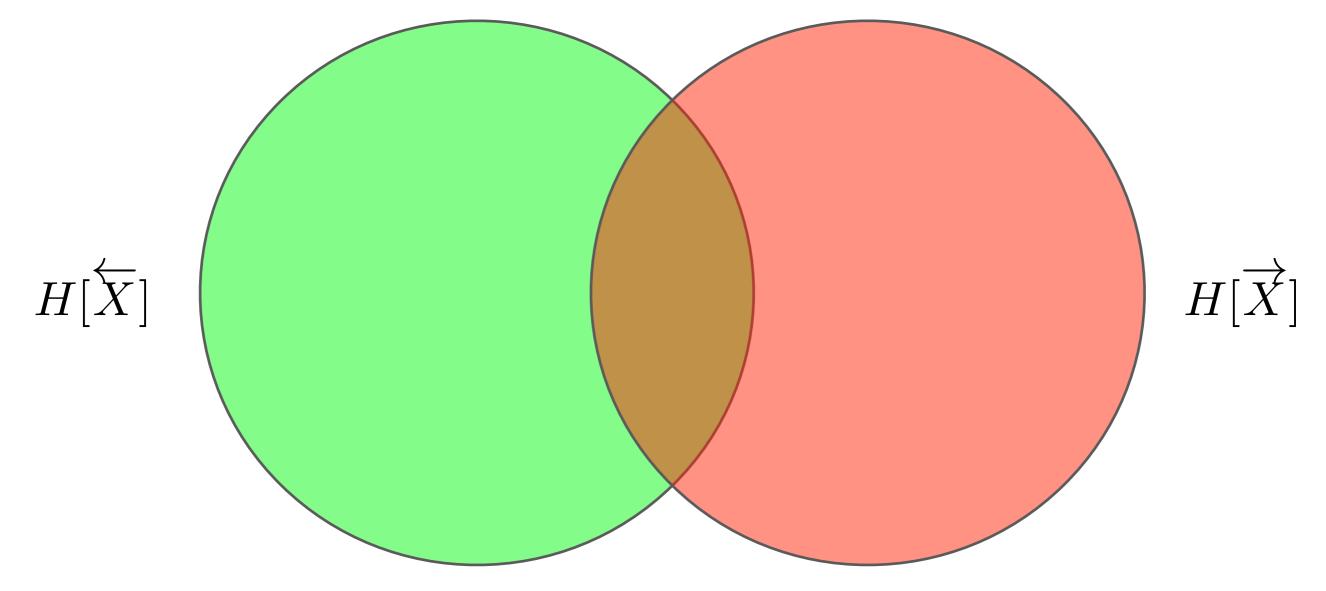
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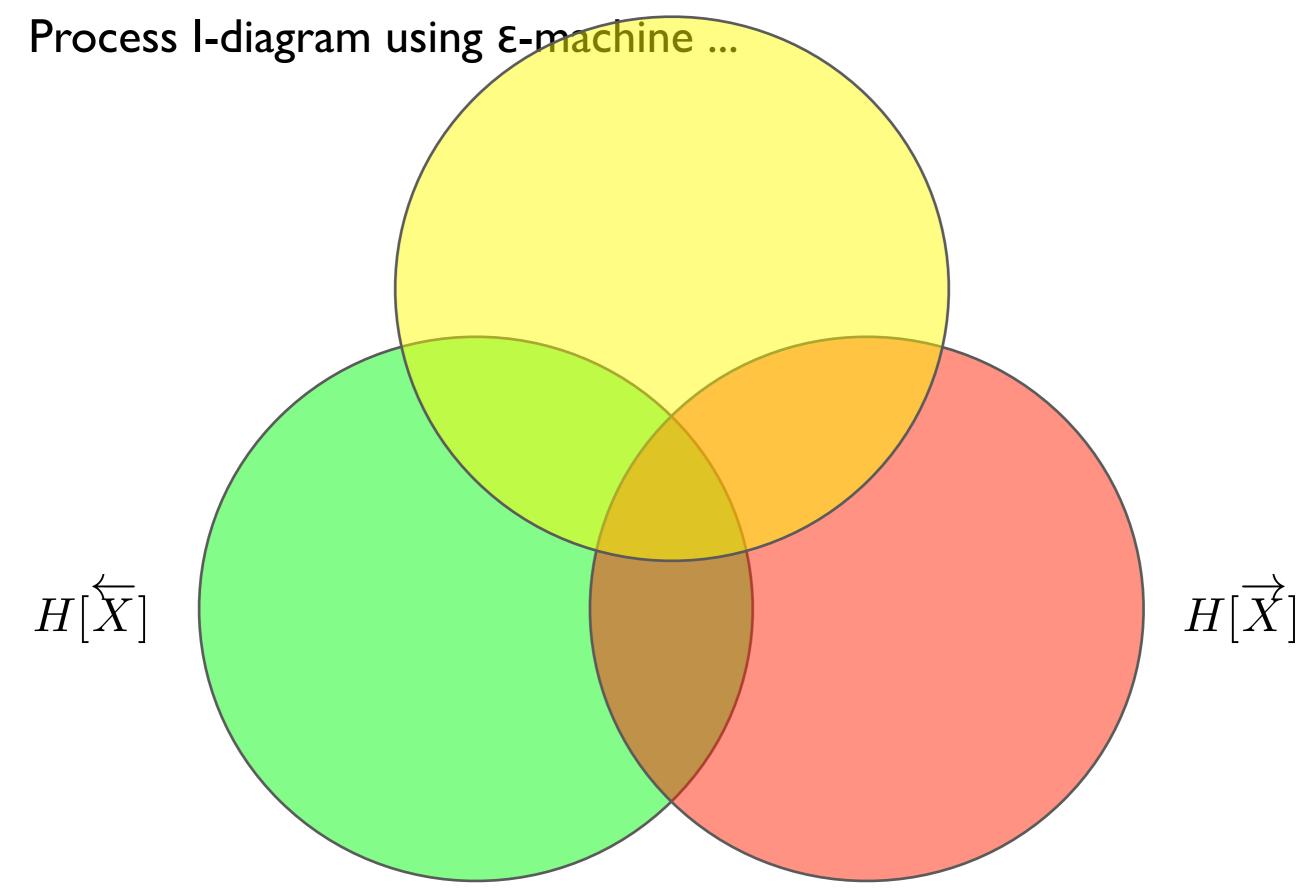
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Information Diagrams for Processes Process I-diagram using E-machine ...

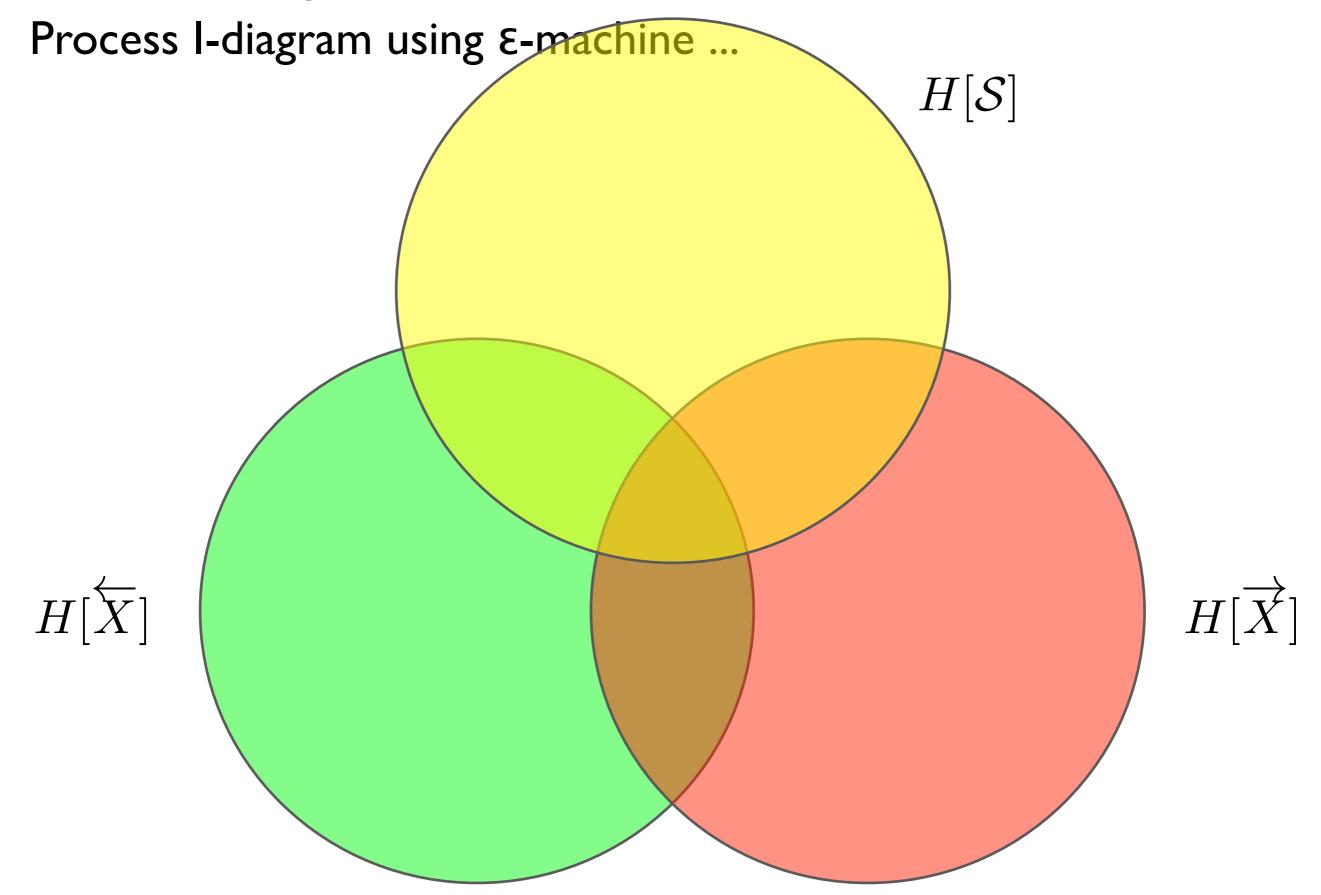


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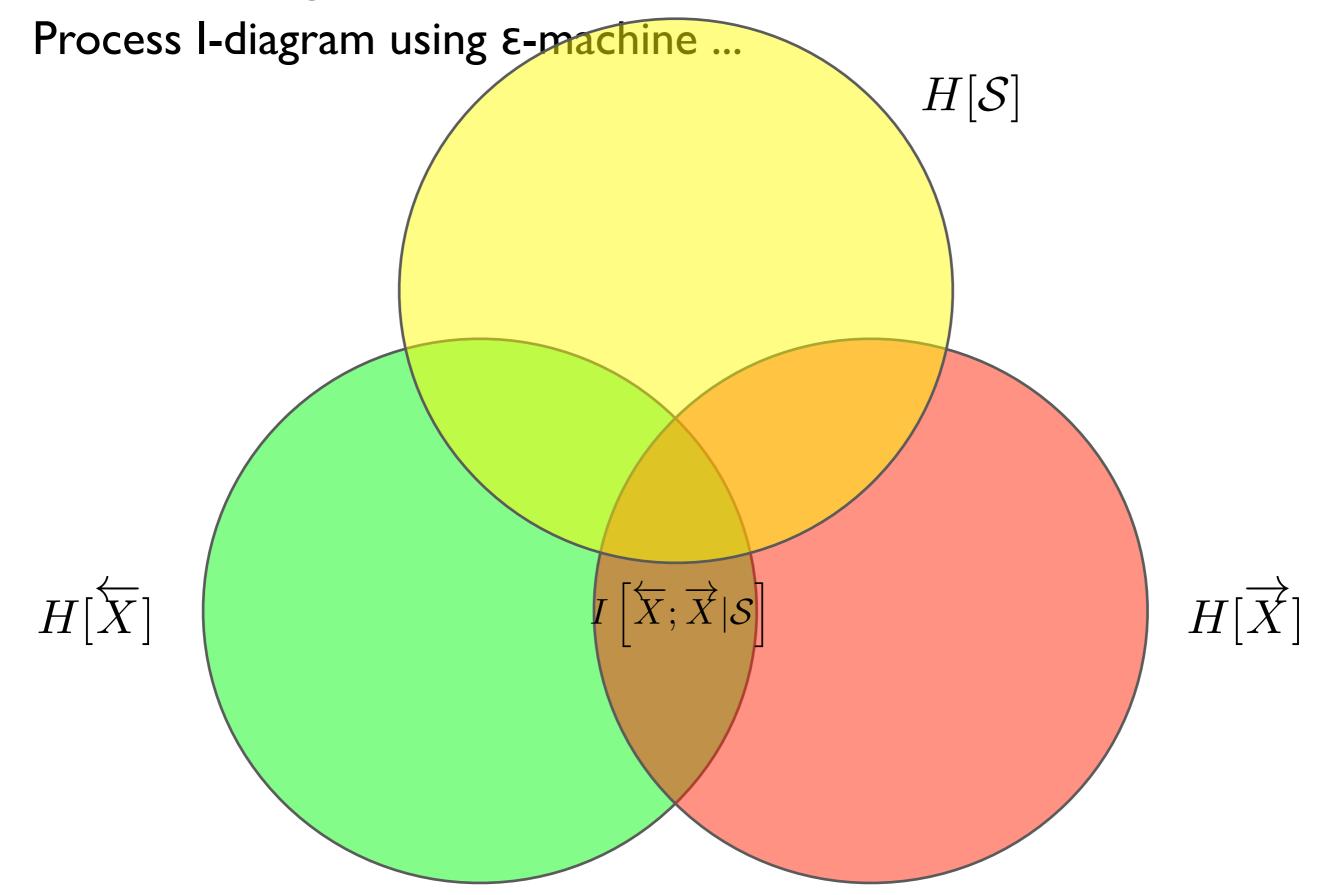




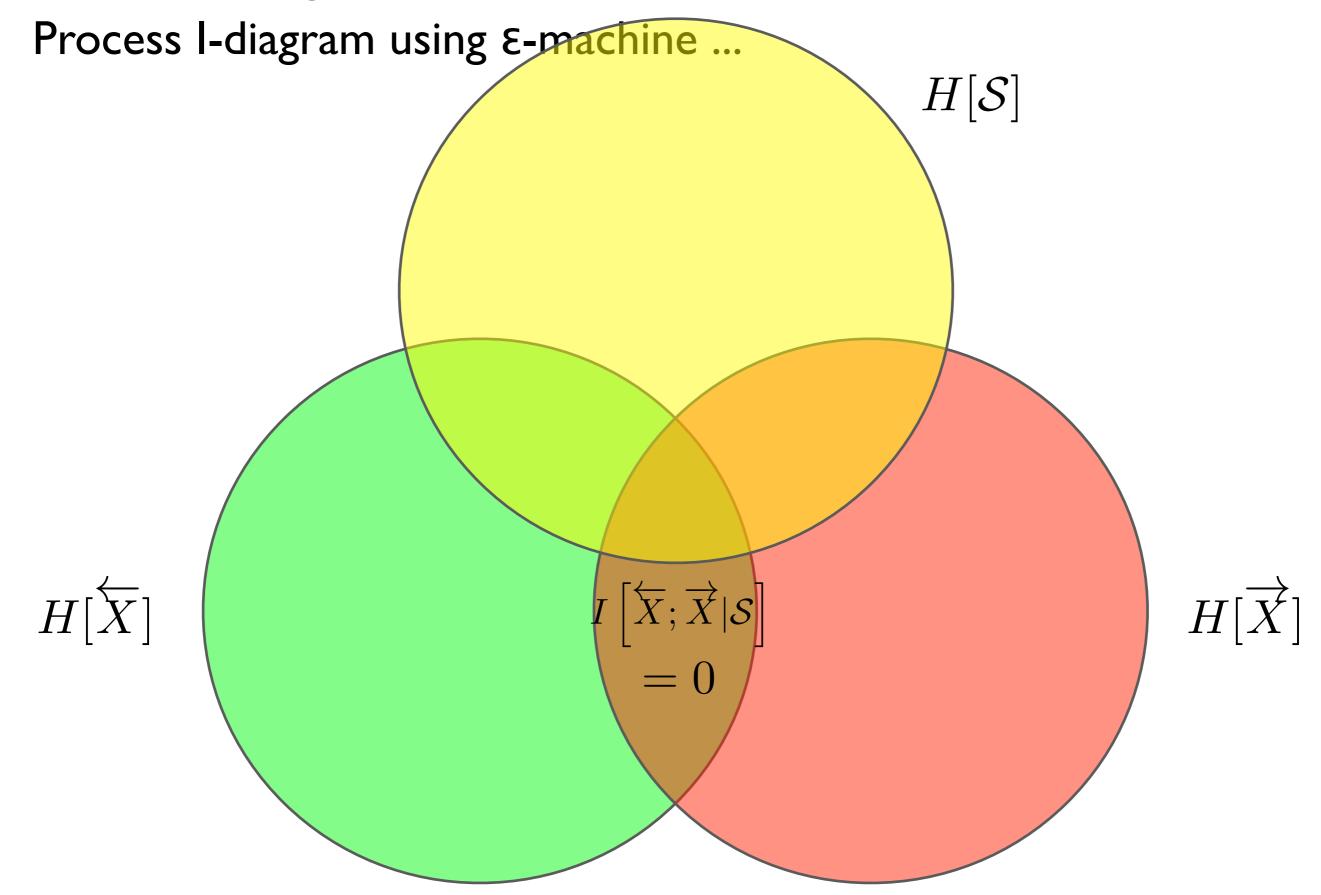
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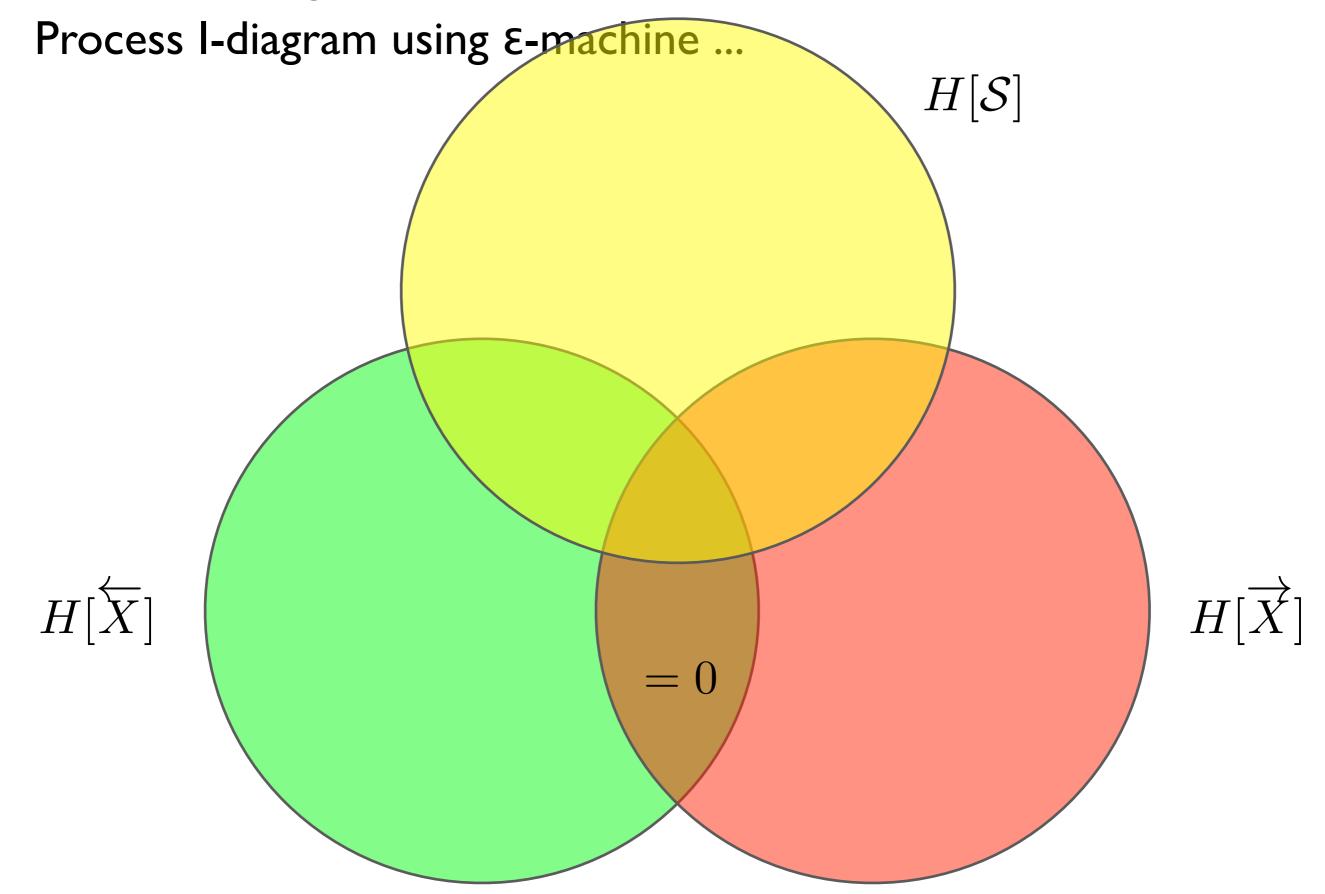
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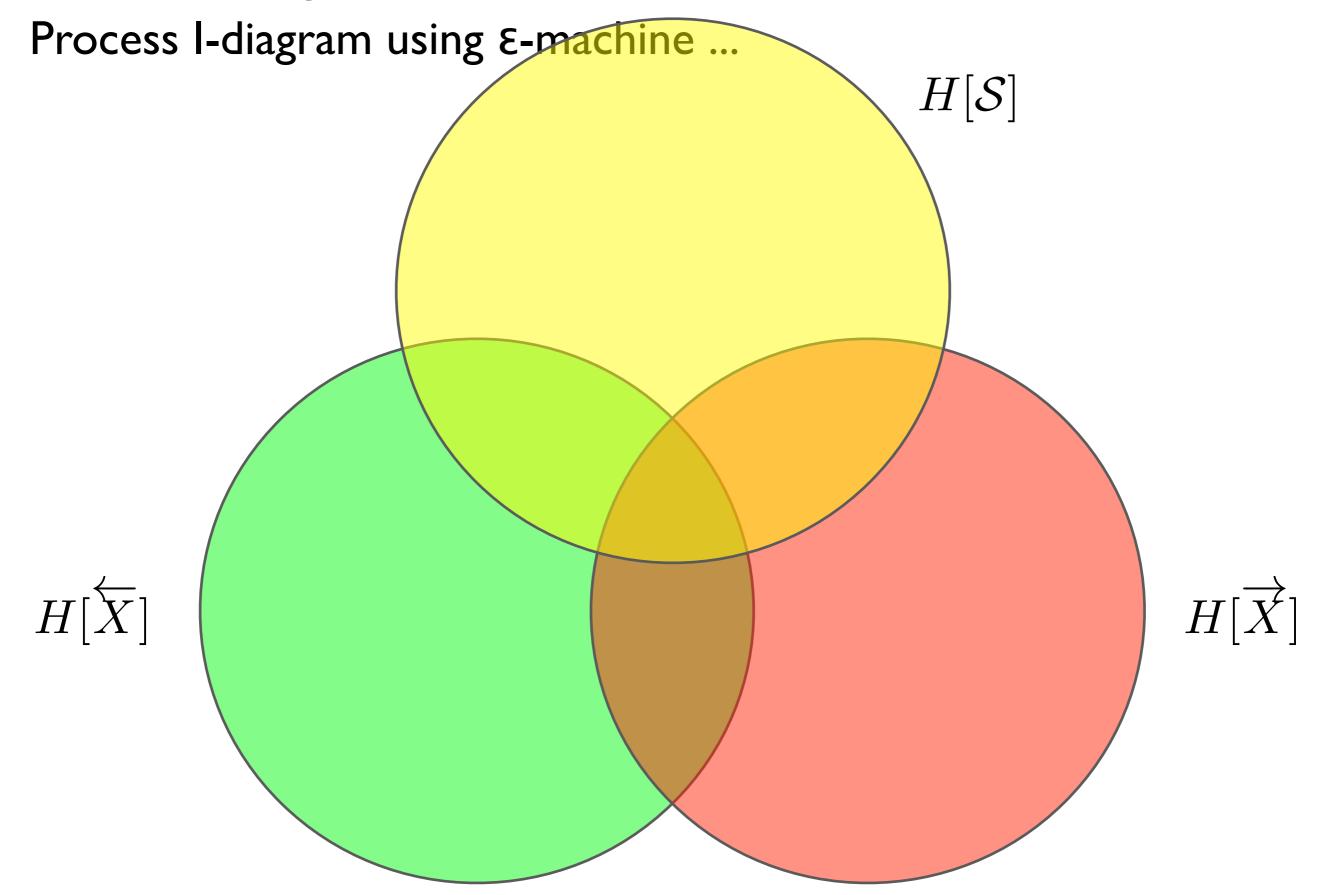
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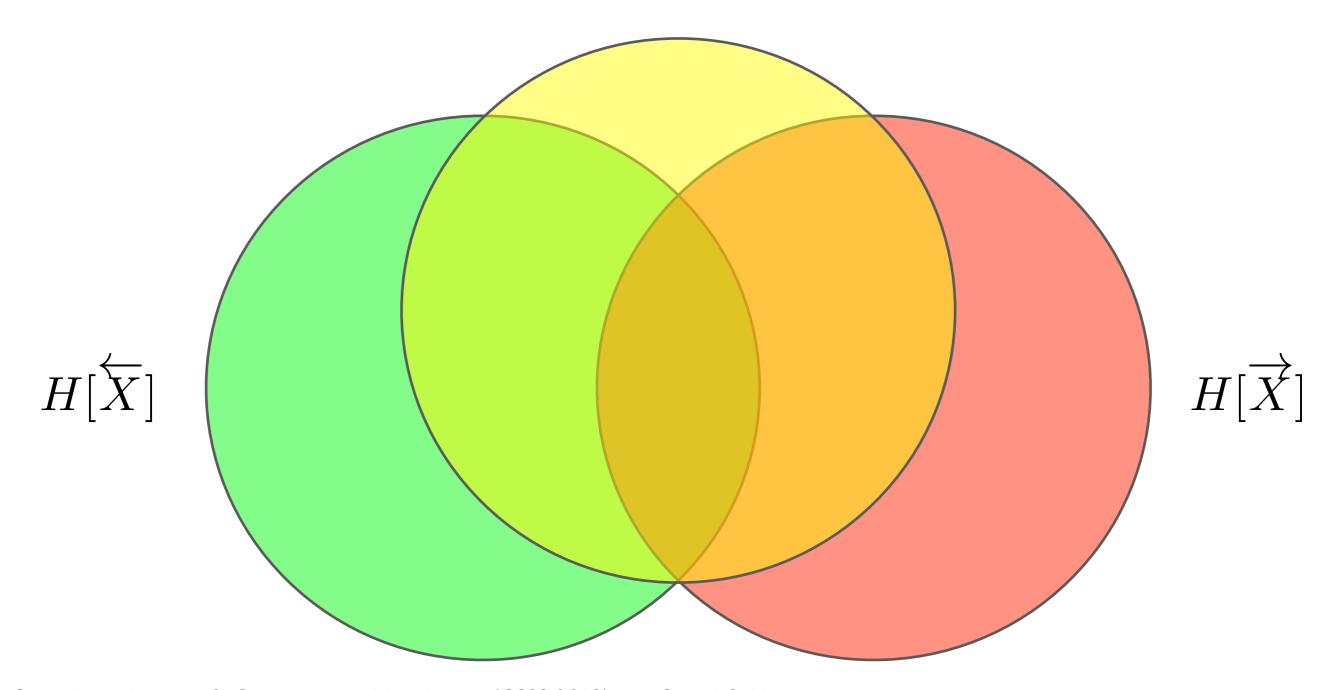
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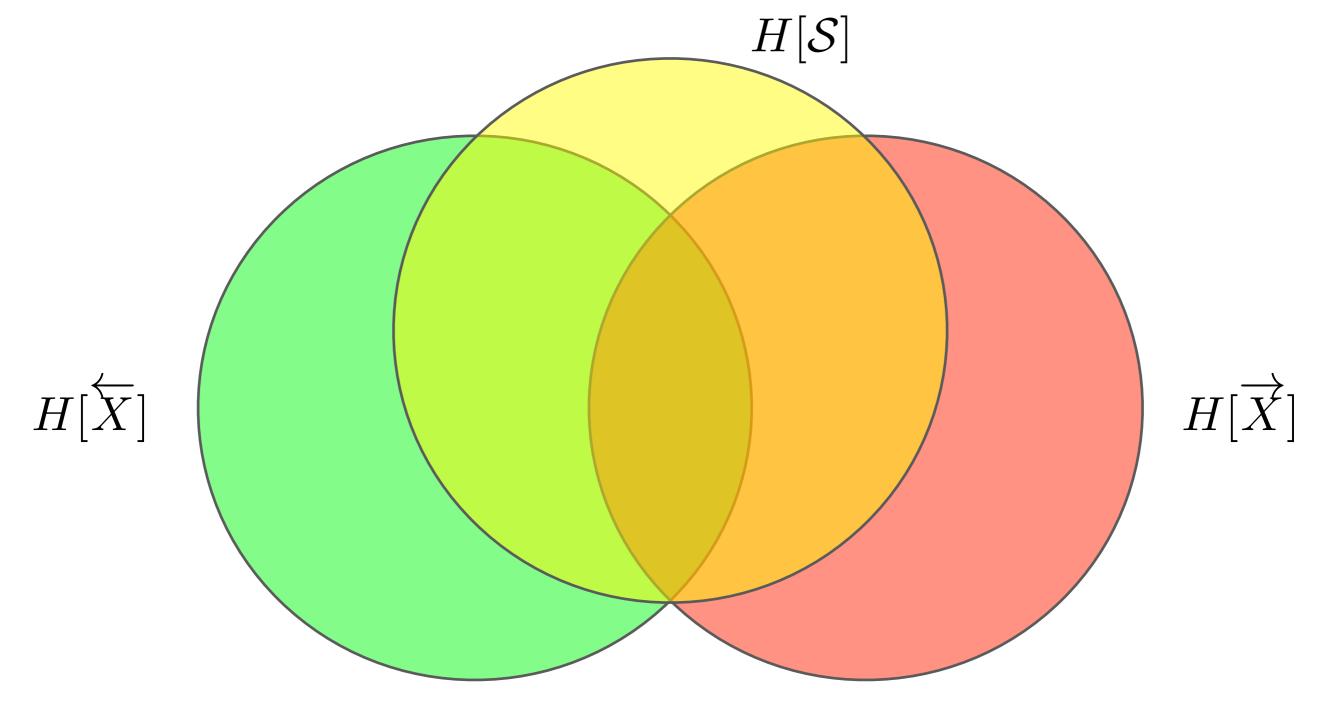
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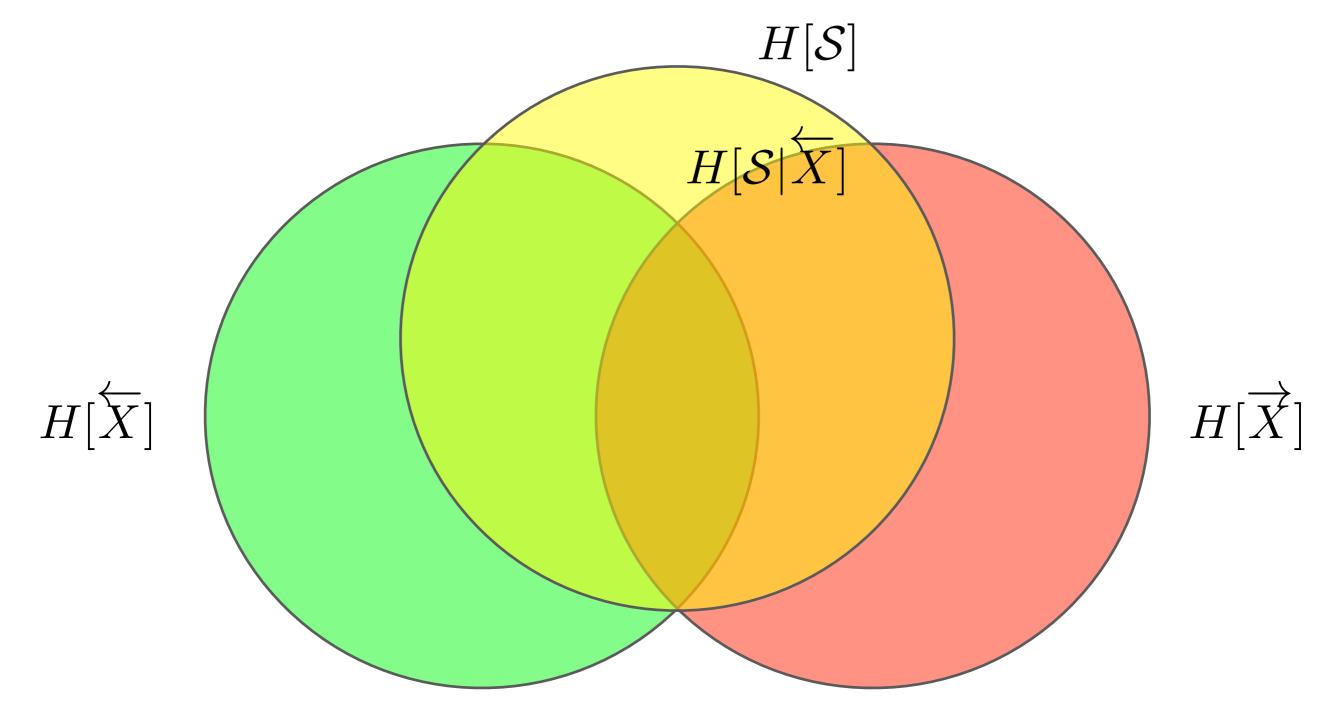
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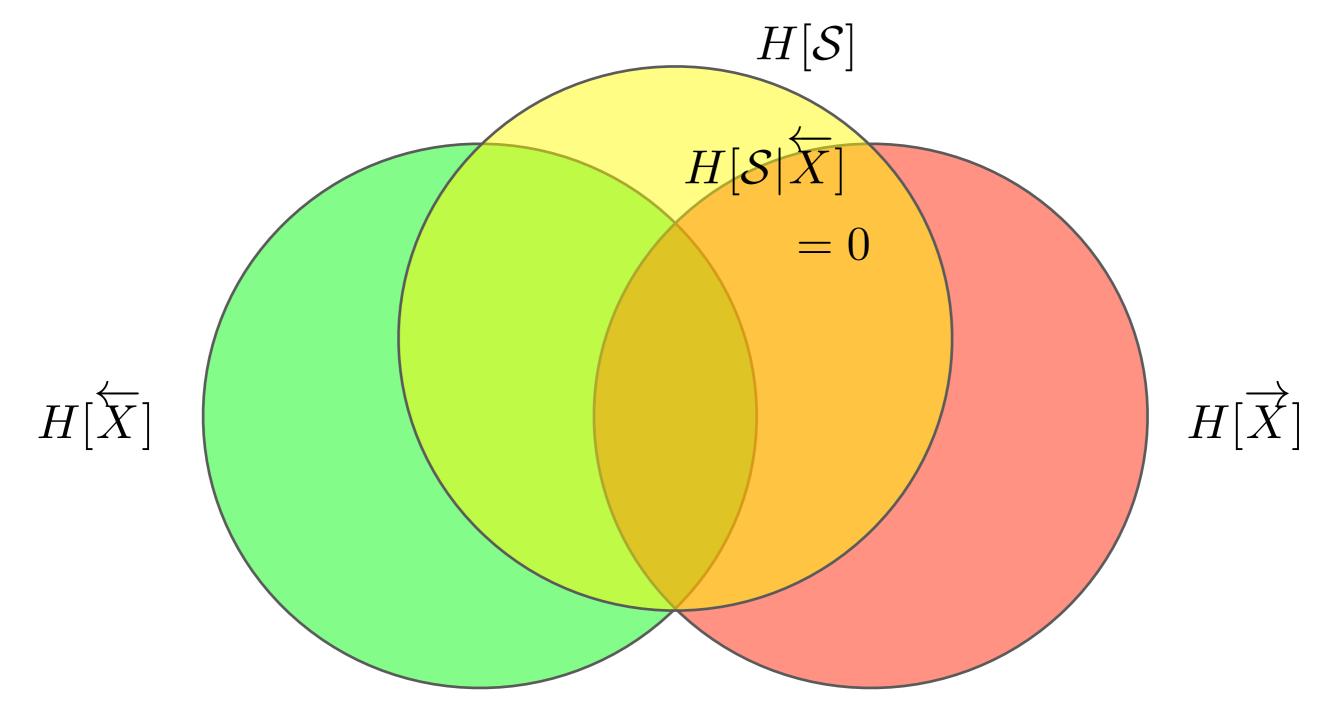
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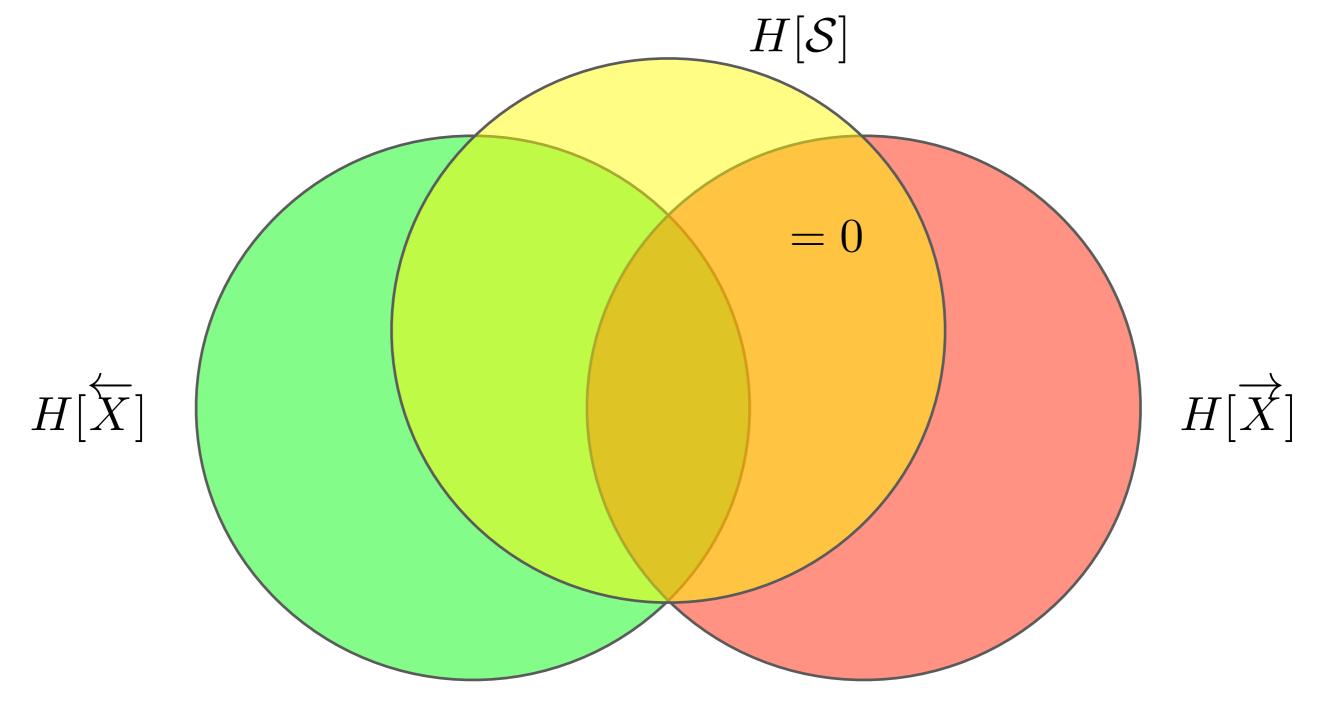
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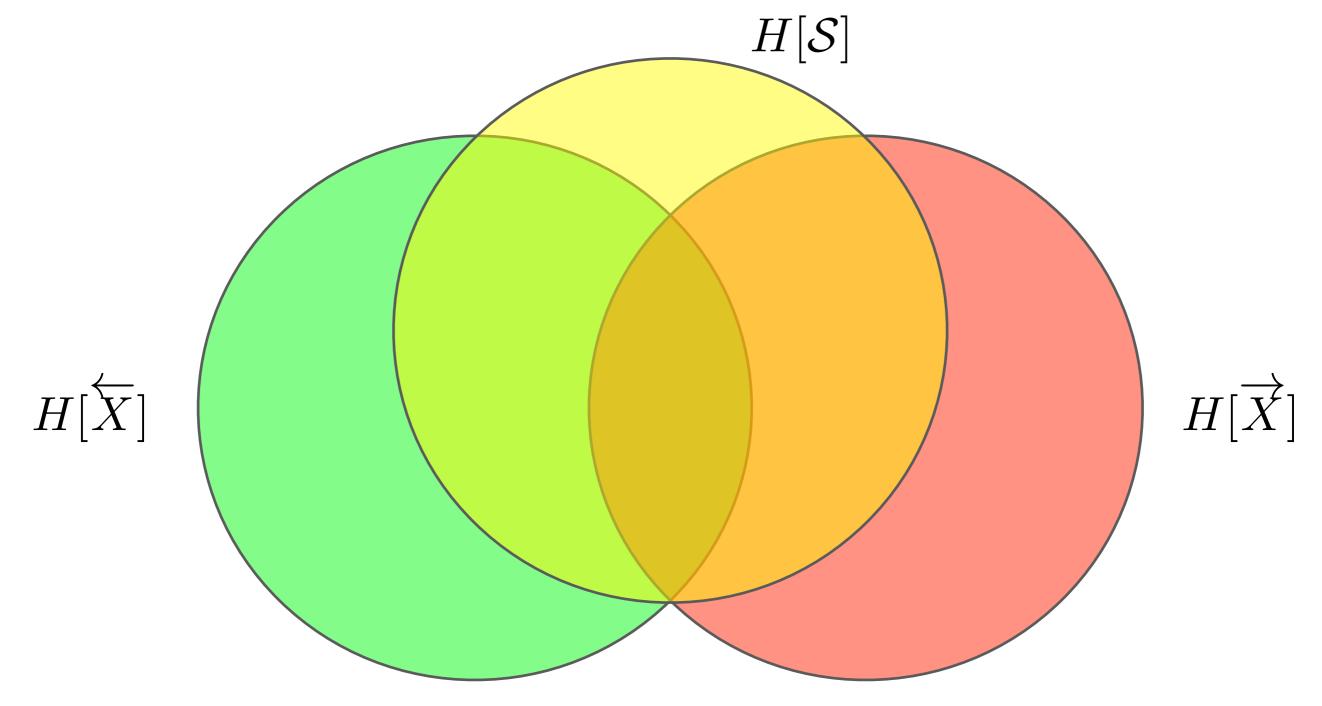
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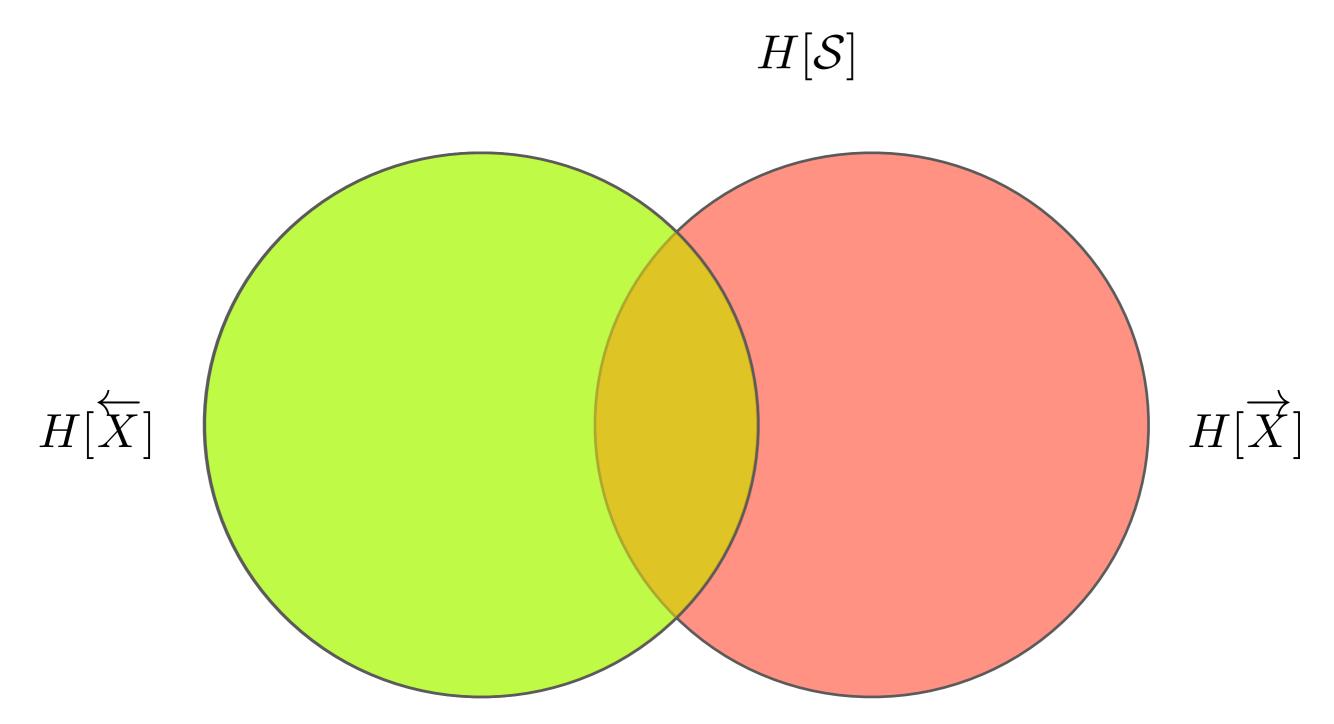
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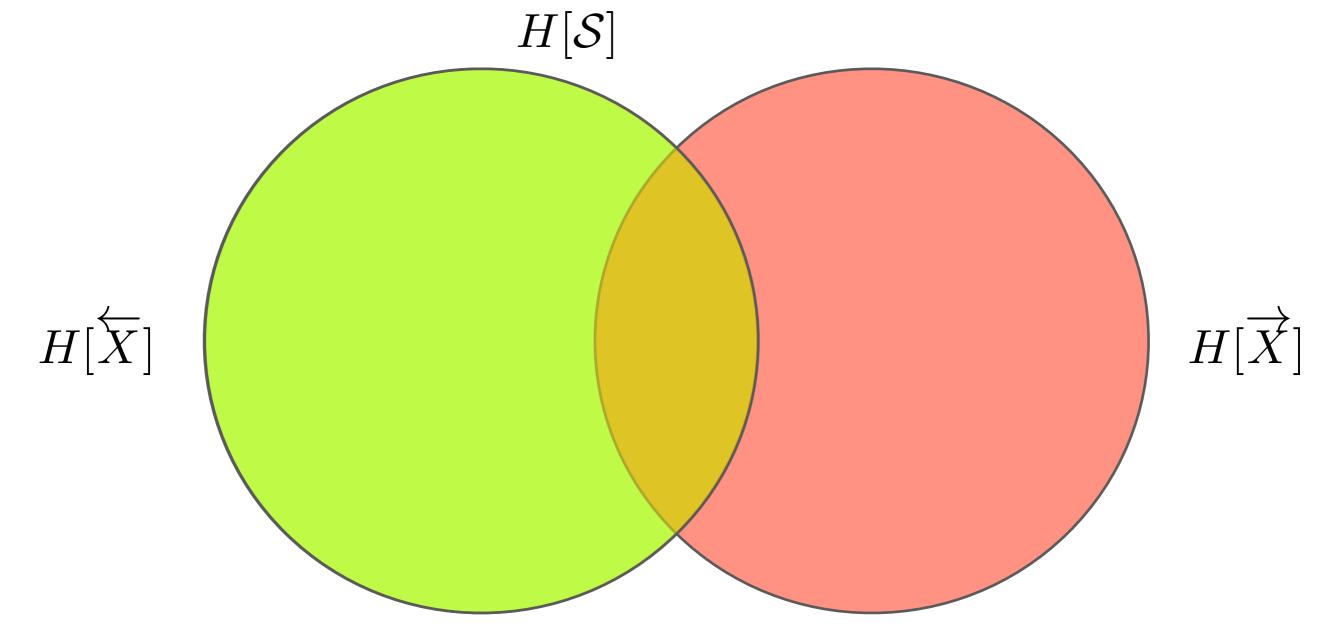
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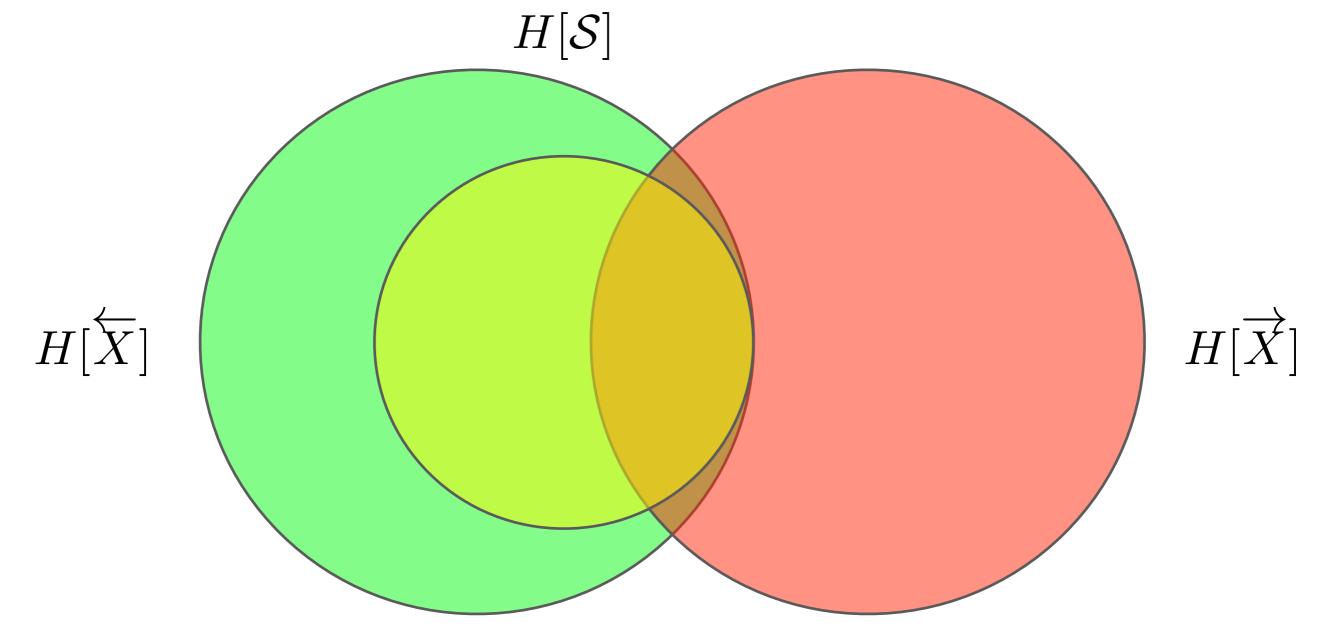
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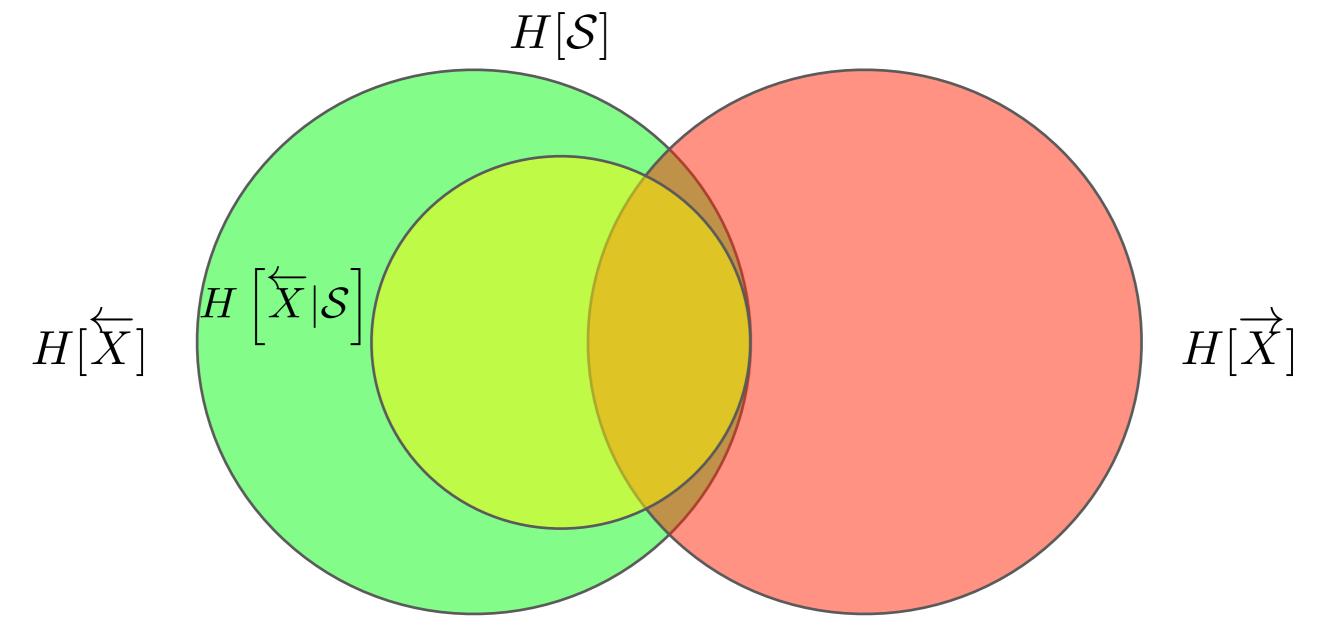
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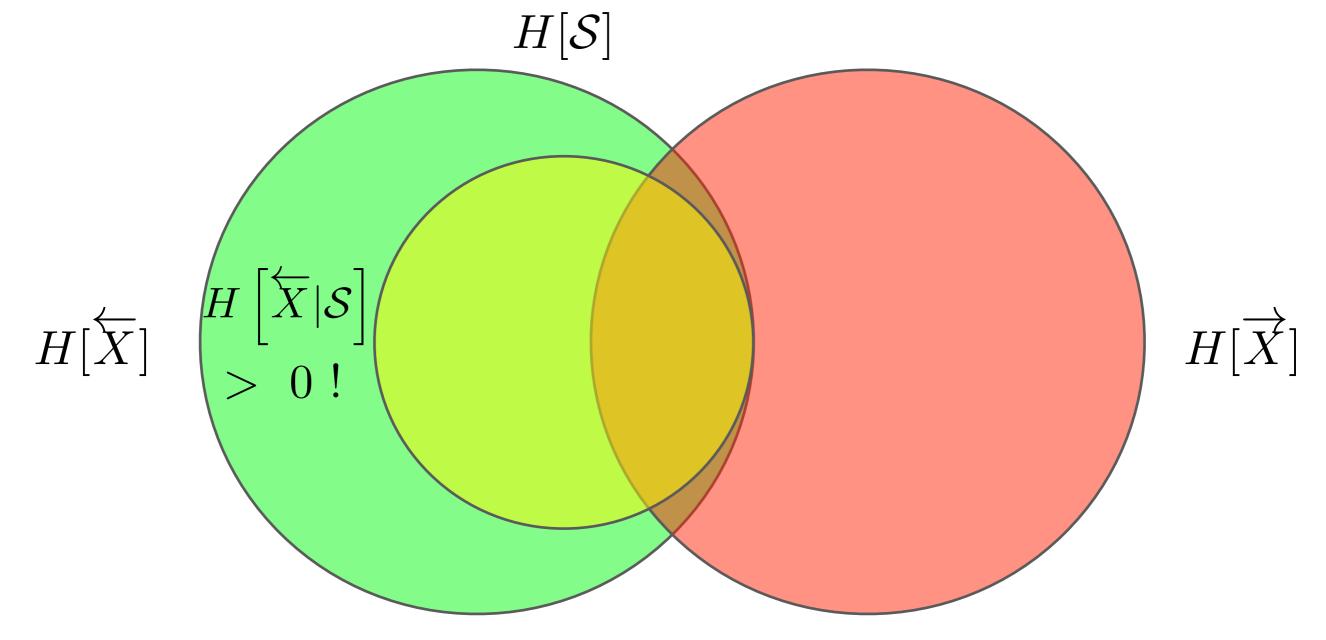
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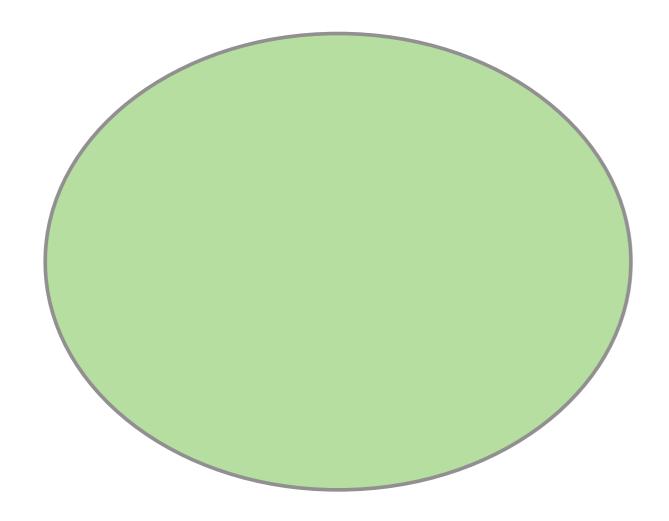


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ε-Machine I-diagram:

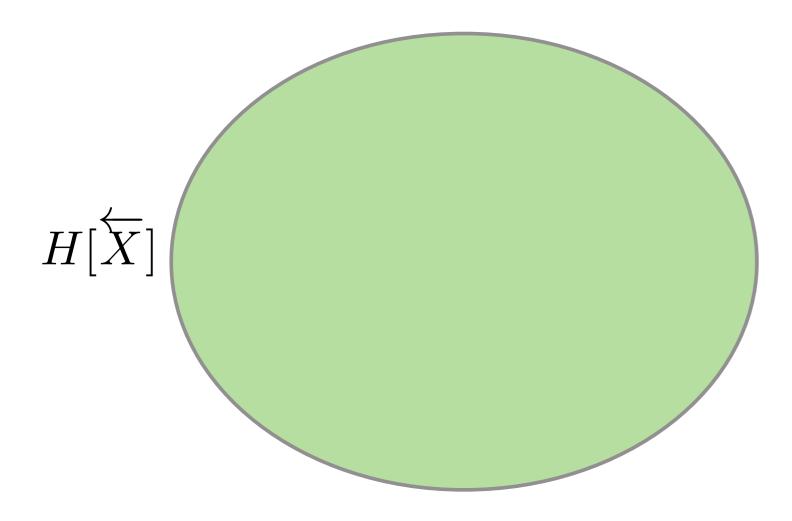
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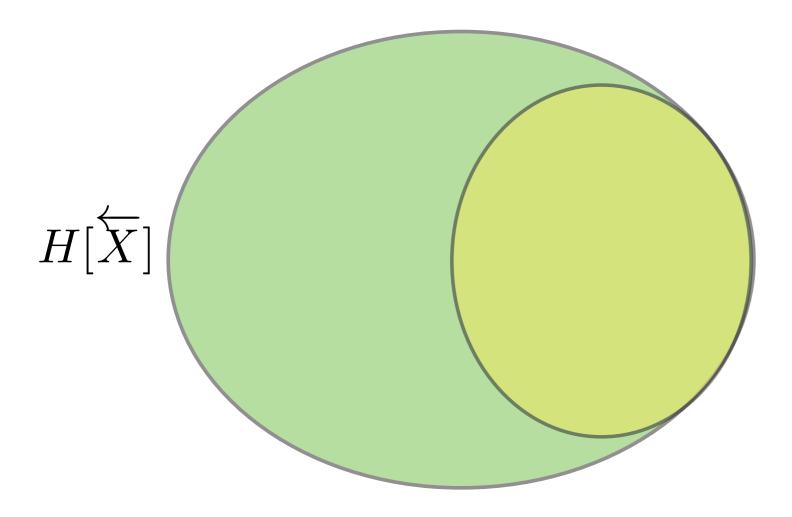
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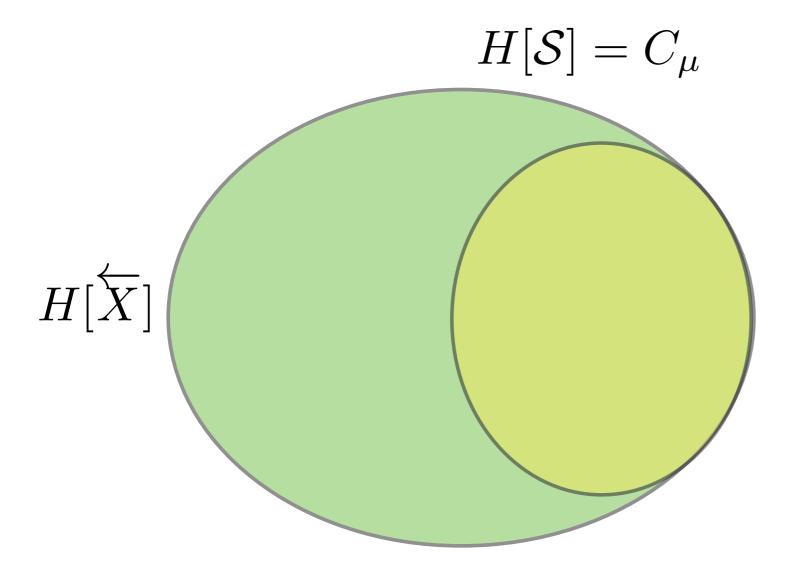
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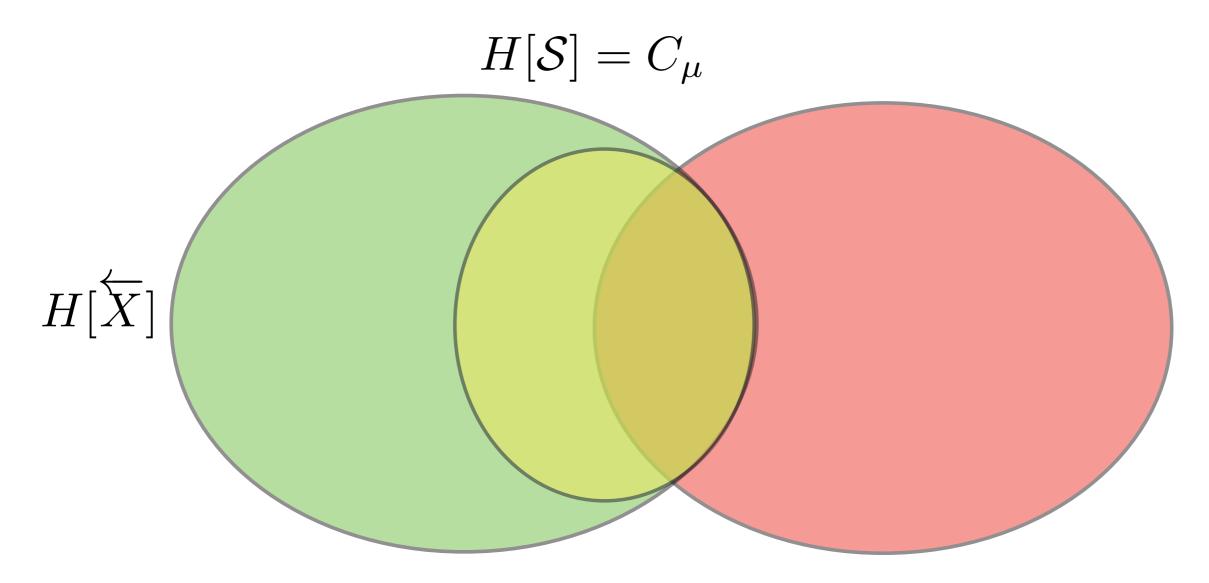
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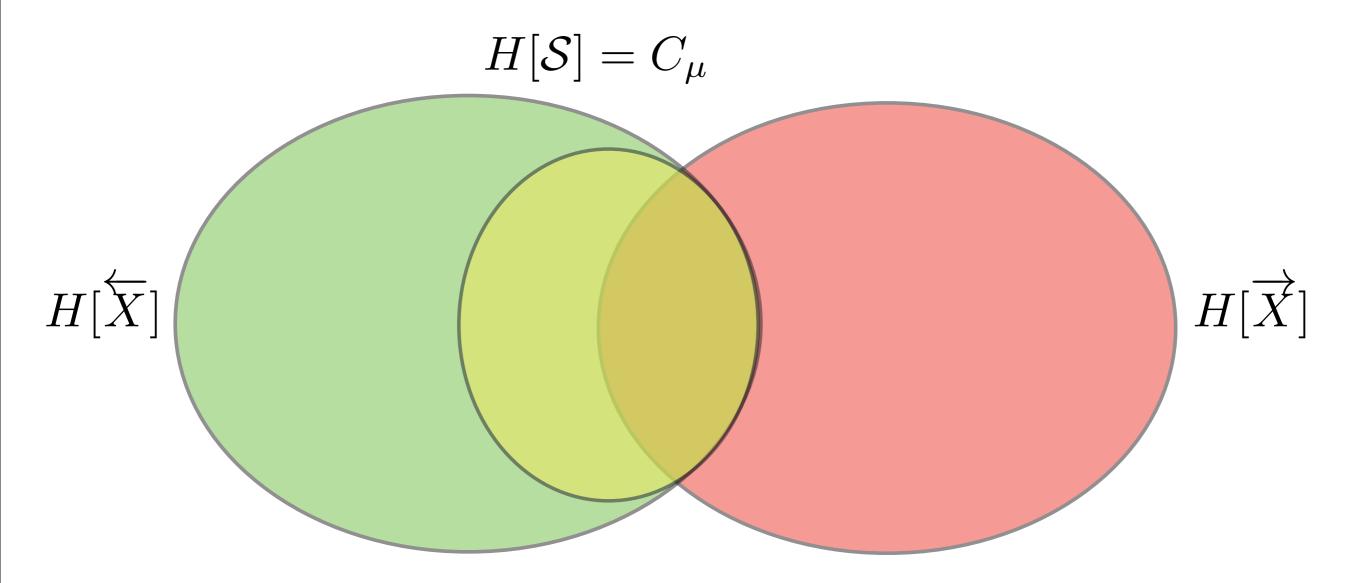
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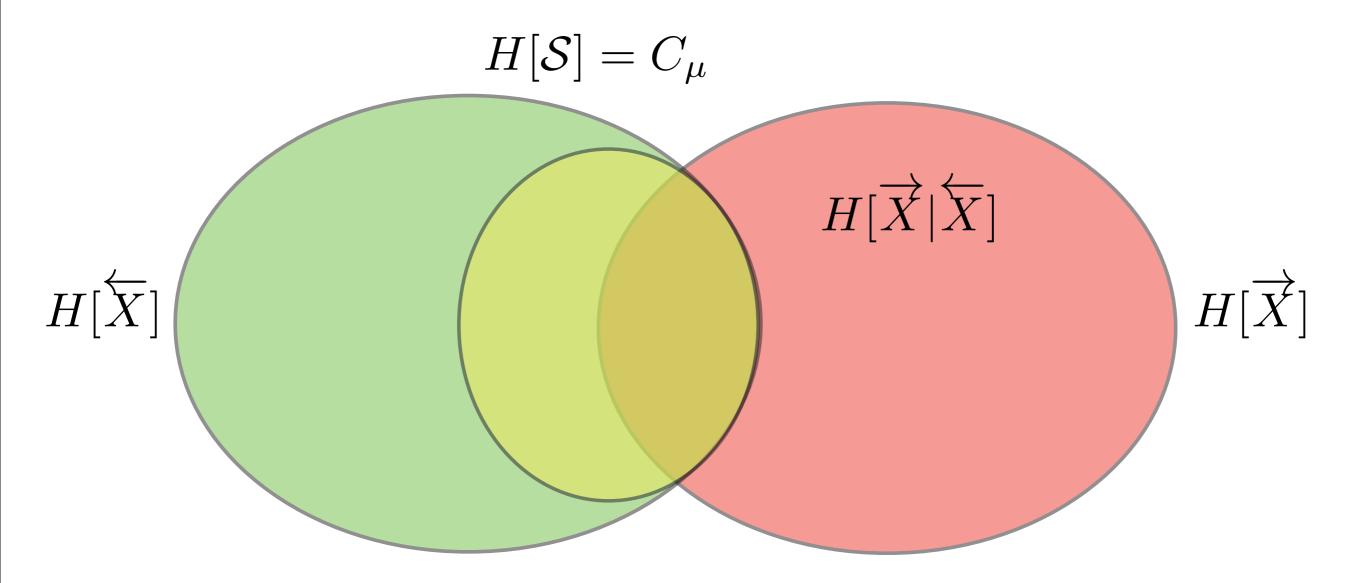
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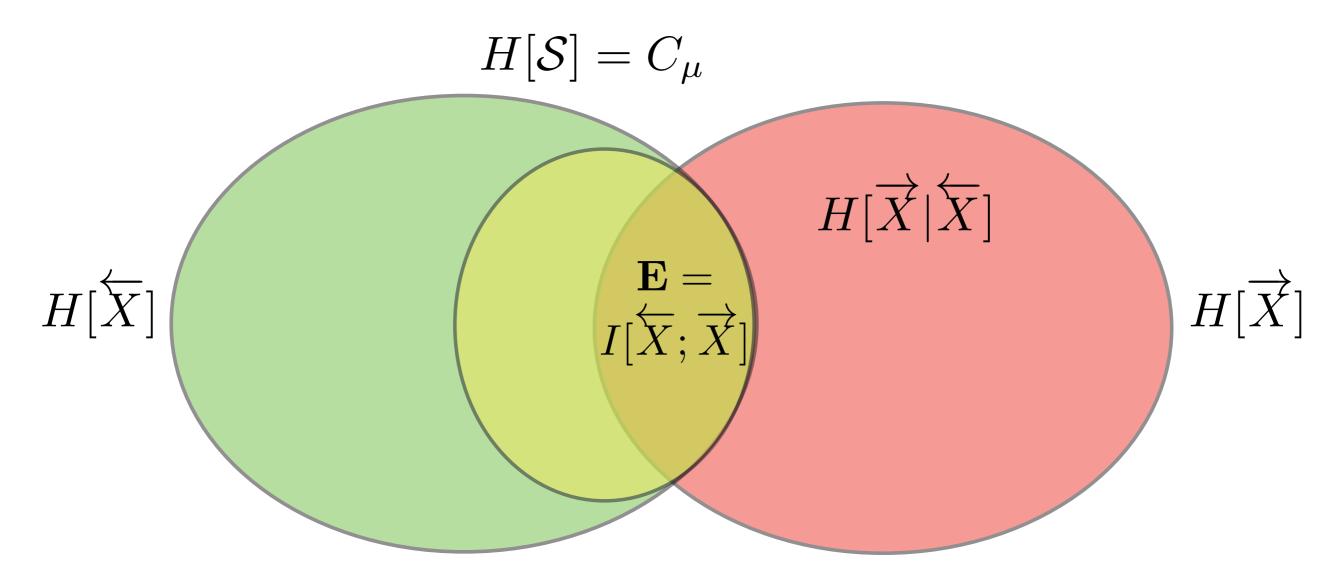
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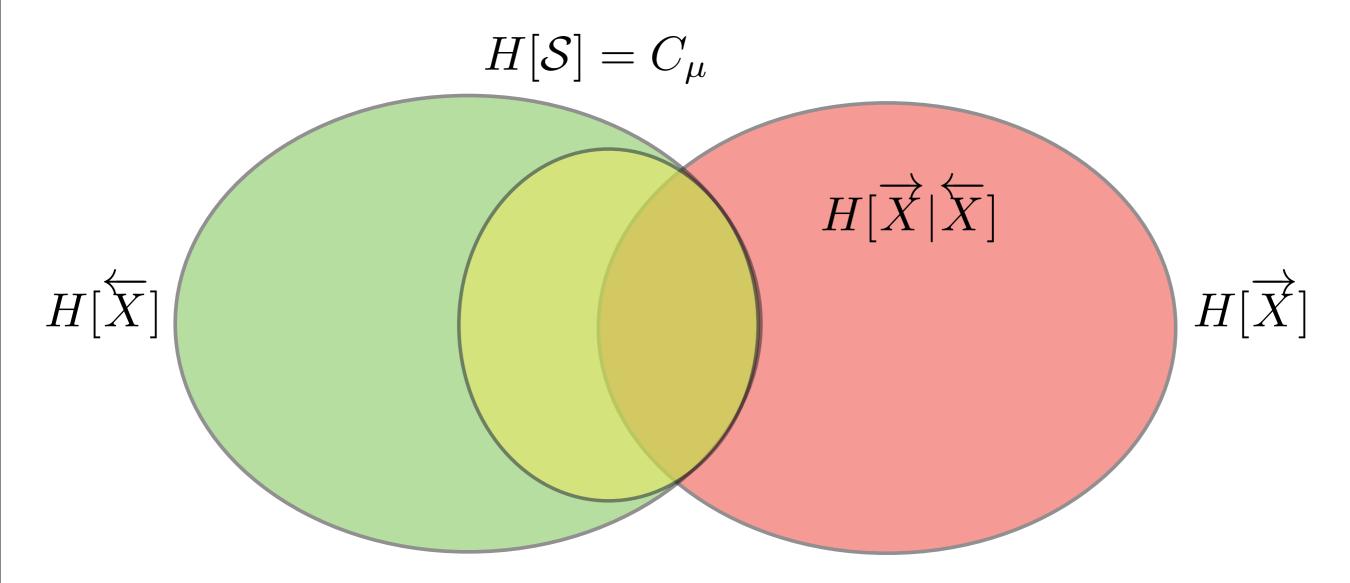
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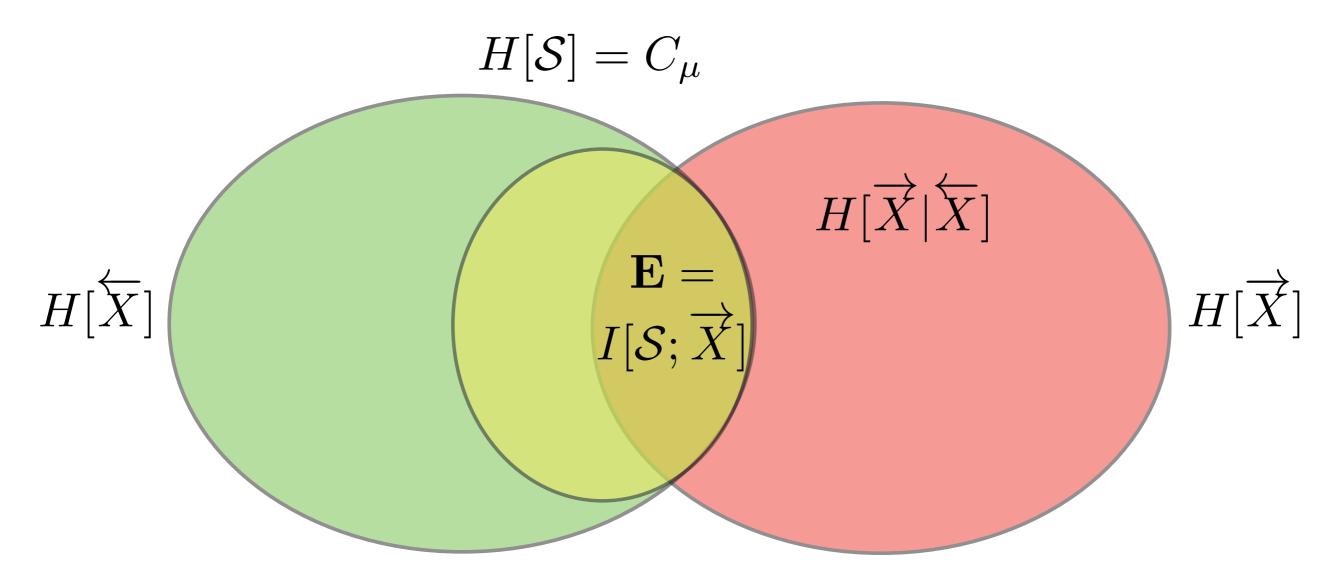
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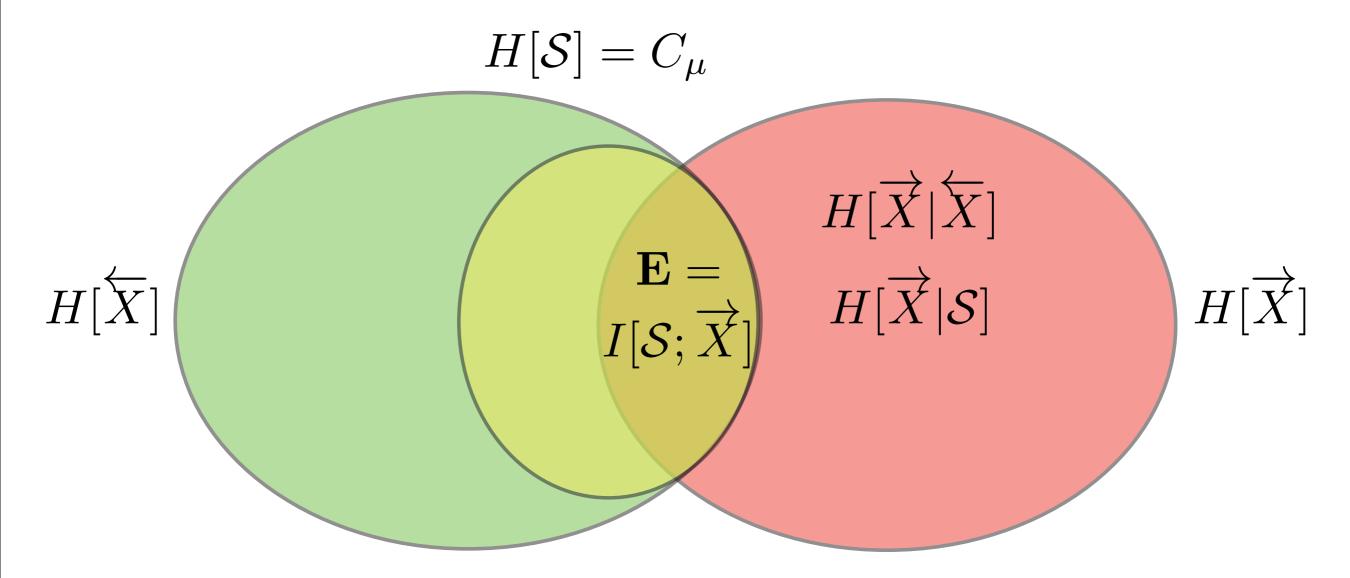
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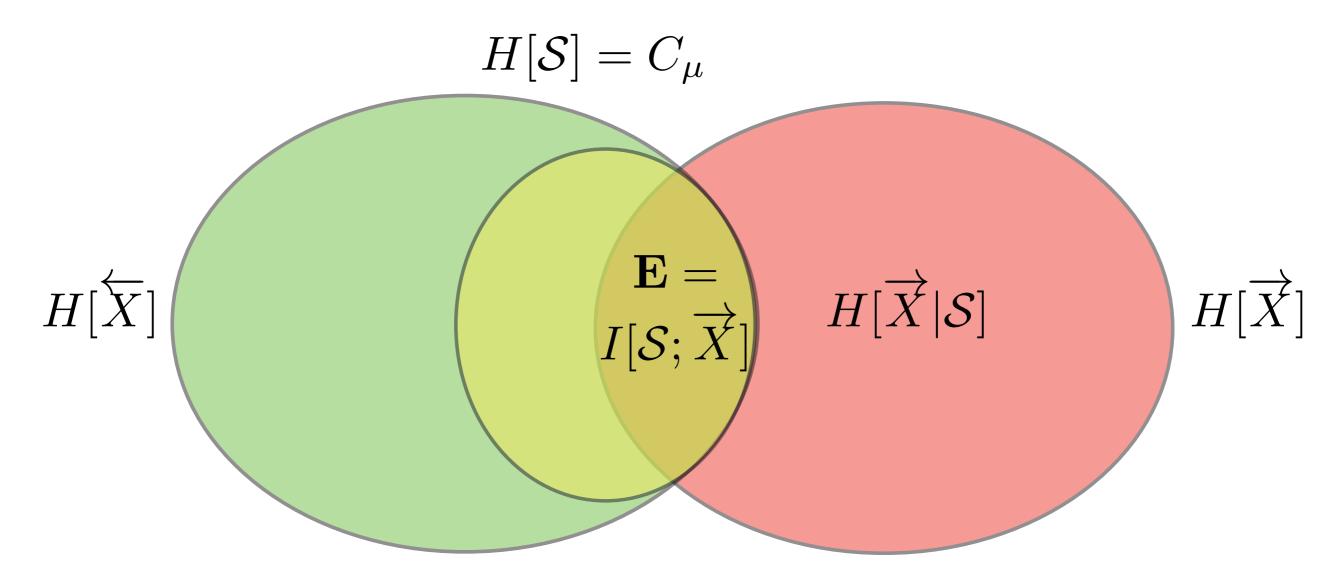
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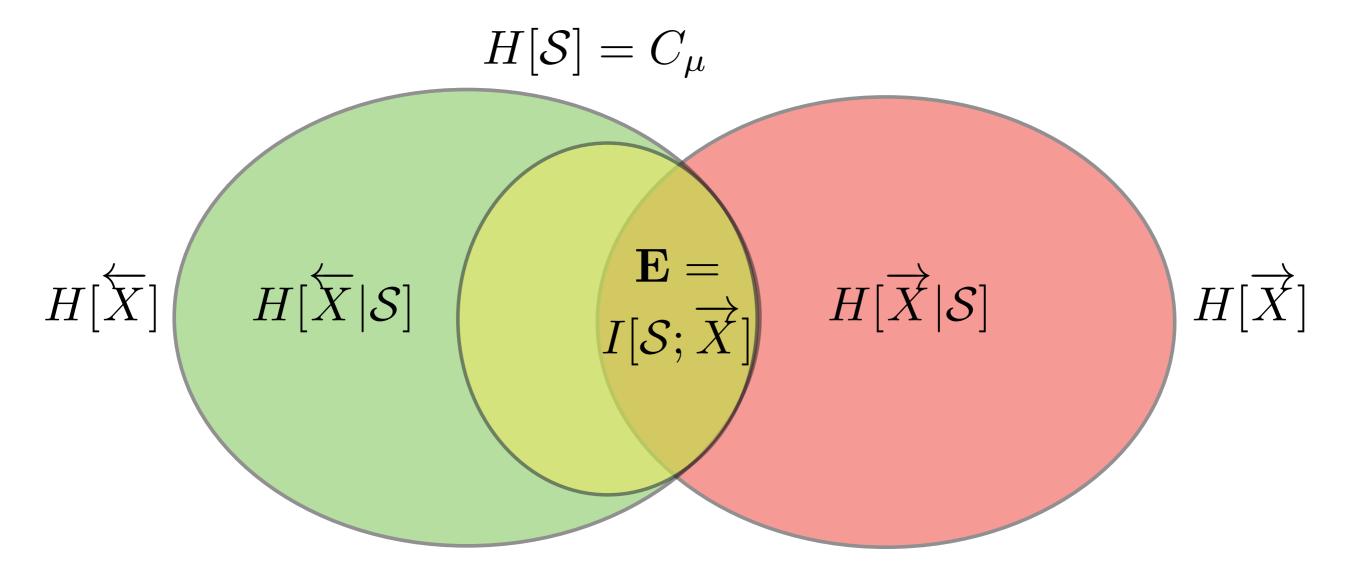
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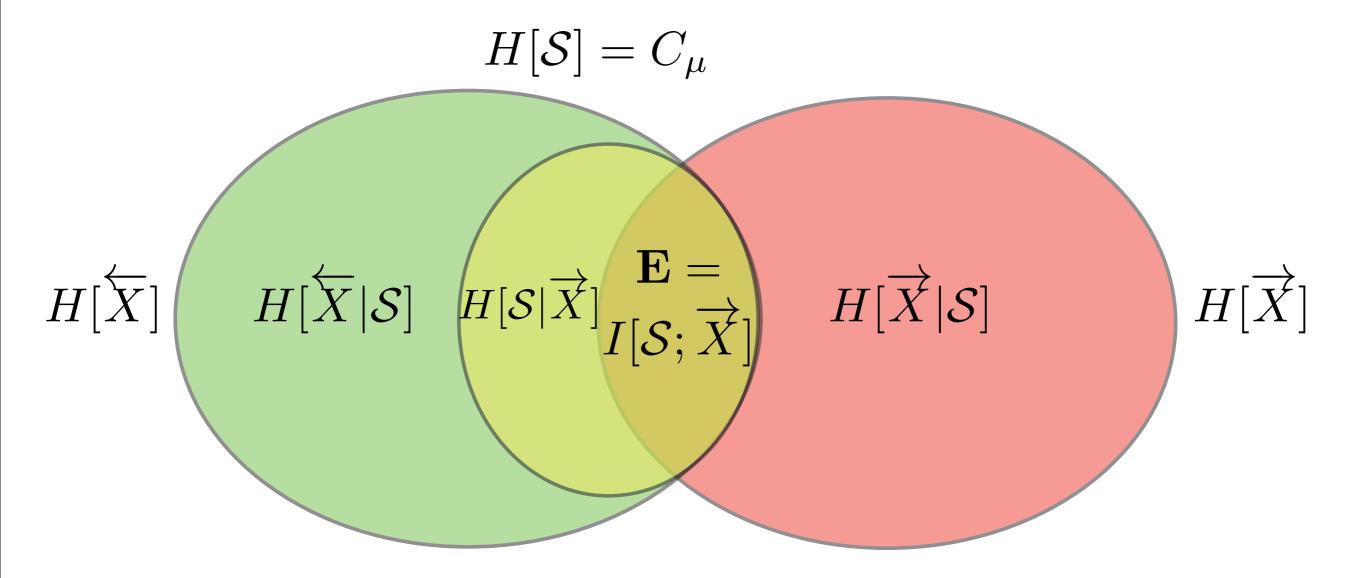
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ε-Machine I-diagram:



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ε-Machine I-diagram:



Complexity Lecture 2: Computational Mechanics (CSSS 2012); Jim Crutchfield

What is $H[\overrightarrow{X}|\mathcal{S}]$?

Unpredictability:
$$H[\overrightarrow{X}^L|\mathcal{S}] = Lh_{\mu}$$

Proof Sketch:

$$H[\overrightarrow{X}^{L}|\mathcal{S}] = H[\overrightarrow{X}^{L}|\overleftarrow{X}]$$

$$= H[X_{0}X_{1} \dots X_{L-1}|\overleftarrow{X}]$$

$$= H[X_{1} \dots X_{L-1}|\overleftarrow{X}X_{0}] + H[X_{0}|\overleftarrow{X}]$$

$$= H[X_{1} \dots X_{L-1}|\overleftarrow{X}] + H[X_{0}|\overleftarrow{X}]$$

$$\vdots$$

$$= H[X_{L-1}|\overleftarrow{X}] + \dots + H[X_{1}|\overleftarrow{X}] + H[X_{0}|\overleftarrow{X}]$$

$$= LH[X_{0}|\overleftarrow{X}]$$

$$= Lh_{\mu}$$

Complexity Lecture 2: Computational Mechanics (CSSS 2012); Jim Crutchfield

Information Diagrams for Processes What is Mystery Wedge? $H[\mathcal{S}|\overrightarrow{X}]$ Uncertainty of causal state given future. Implications?

Recall Bound on Excess Entropy: $\mathbf{E} \leq C_{\mu}$

Proof sketch:
$$\mathbf{E} = I[\overleftarrow{X}; \overrightarrow{X}]$$

$$= H[\overrightarrow{X}] - H[\overrightarrow{X}|\overleftarrow{X}]$$

$$= H[\overrightarrow{X}] - H[\overrightarrow{X}|\mathcal{S}]$$

$$= I[\overrightarrow{X}; \mathcal{S}]$$

$$= H[\mathcal{S}] - H[\mathcal{S}|\overrightarrow{X}]$$

$$\leq H[\mathcal{S}]$$

$$= C_{\mu}$$

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$$= H[S] - H[S|\overrightarrow{X}]$$

$$\leq H[S]$$

$$= C_{\mu}$$

Mystery Wedge!



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$$= H[\mathcal{S}] - H[\mathcal{S}|\overrightarrow{X}]$$

$$< H[\mathcal{S}]$$

Wedge is the inaccessibility of hidden state information!

 $= C_{\prime\prime}$

$$H[\mathcal{S}|\overrightarrow{X}] = C_{\mu} - \mathbf{E}$$
 Wedge controls Internal - Observed

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Information Diagrams for Processes What is Mystery Wedge? $H[\mathcal{S}|\overrightarrow{X}]$

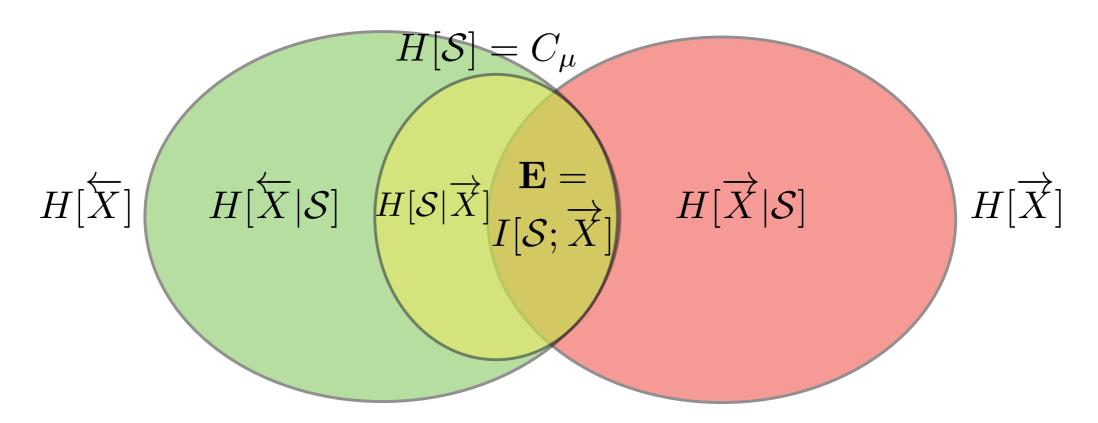
Wedge is the inaccessibility of hidden state information!

$$H[\mathcal{S}|\overrightarrow{X}] = C_{\mu} - \mathbf{E}$$

The process crypticity:

$$\chi = C_{\mu} - \mathbf{E}$$

Controls how much internal state information is observable.



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How to get \mathbf{E} from $\epsilon \mathbf{M}$?

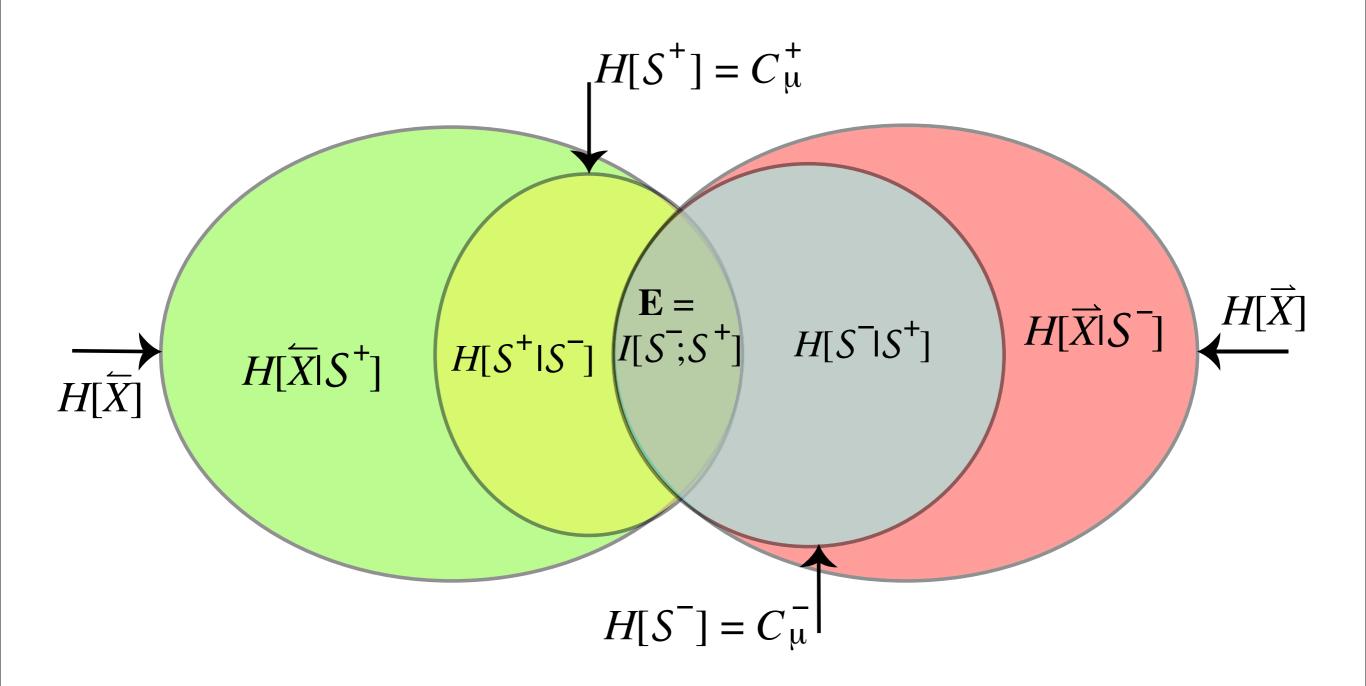
DIRECTIONAL COMPUTATIONAL MECHANICS

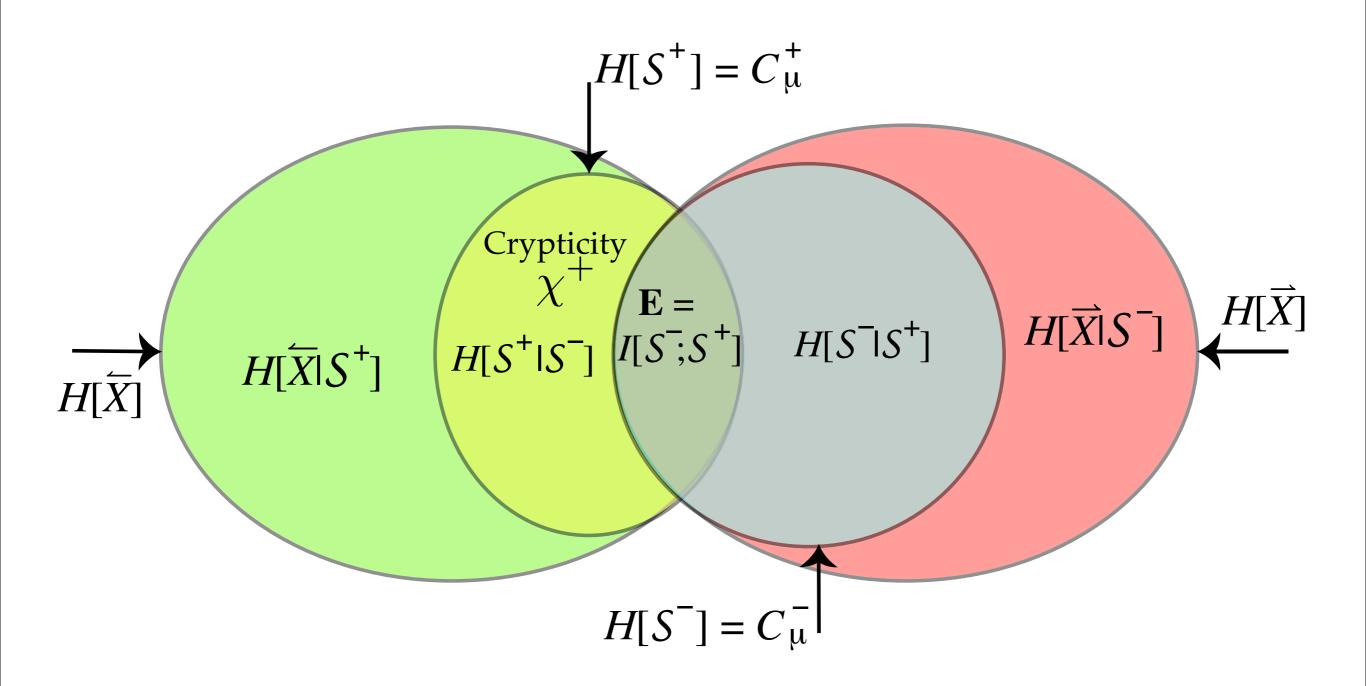
Theorem:

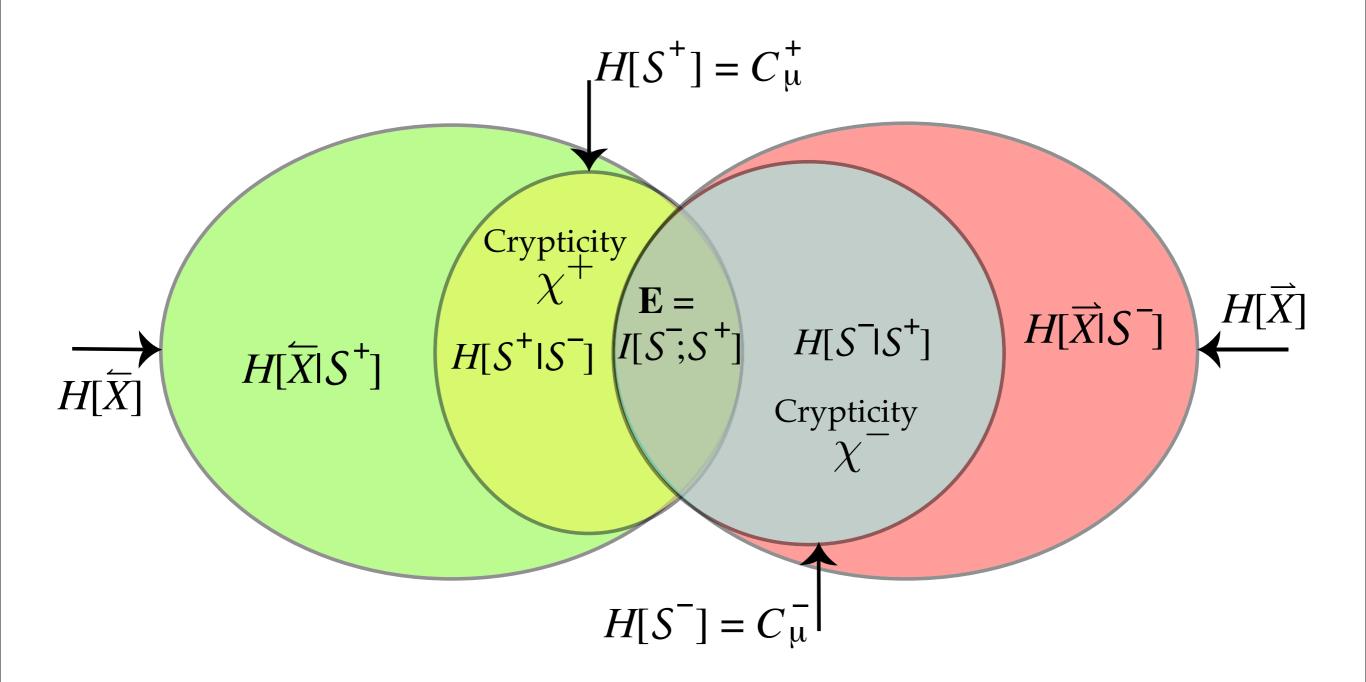
$$\mathbf{E} = I[\mathcal{S}^+; \mathcal{S}^-]$$

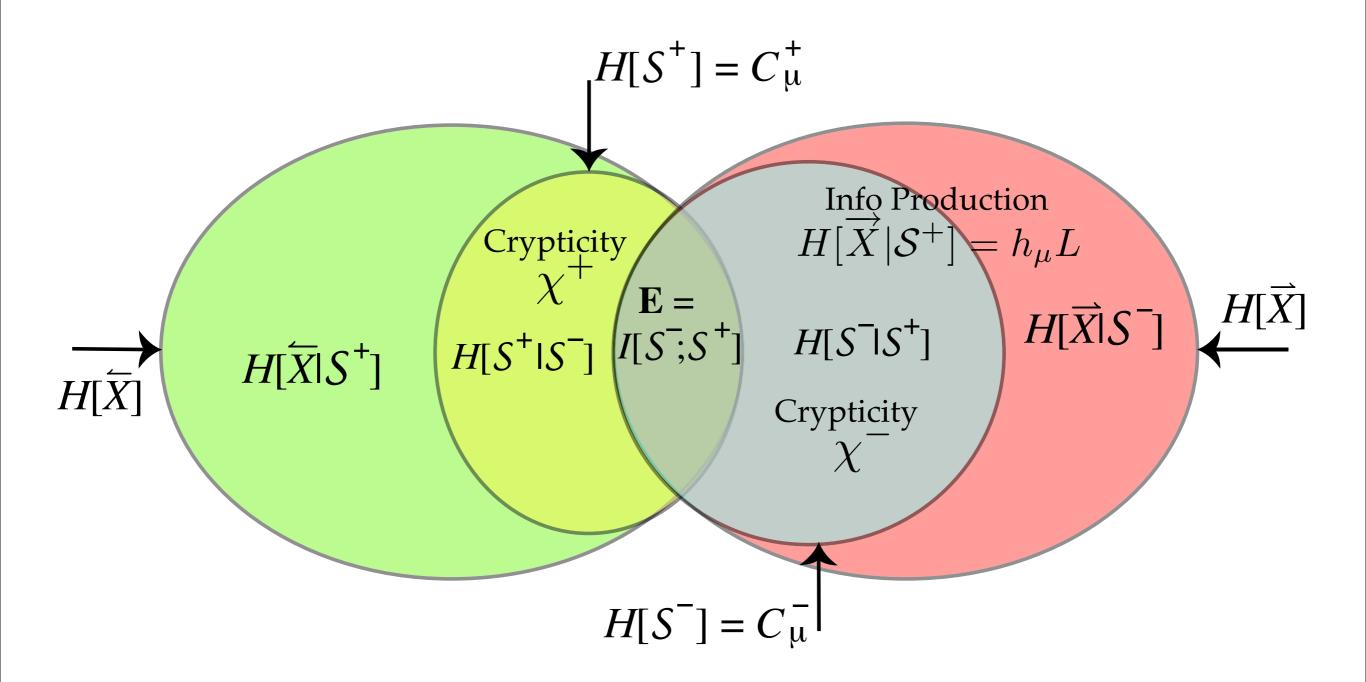
- Effective transmission capacity of channel between forward and reverse processes.
- Time agnostic representation: The **BiMachine**.

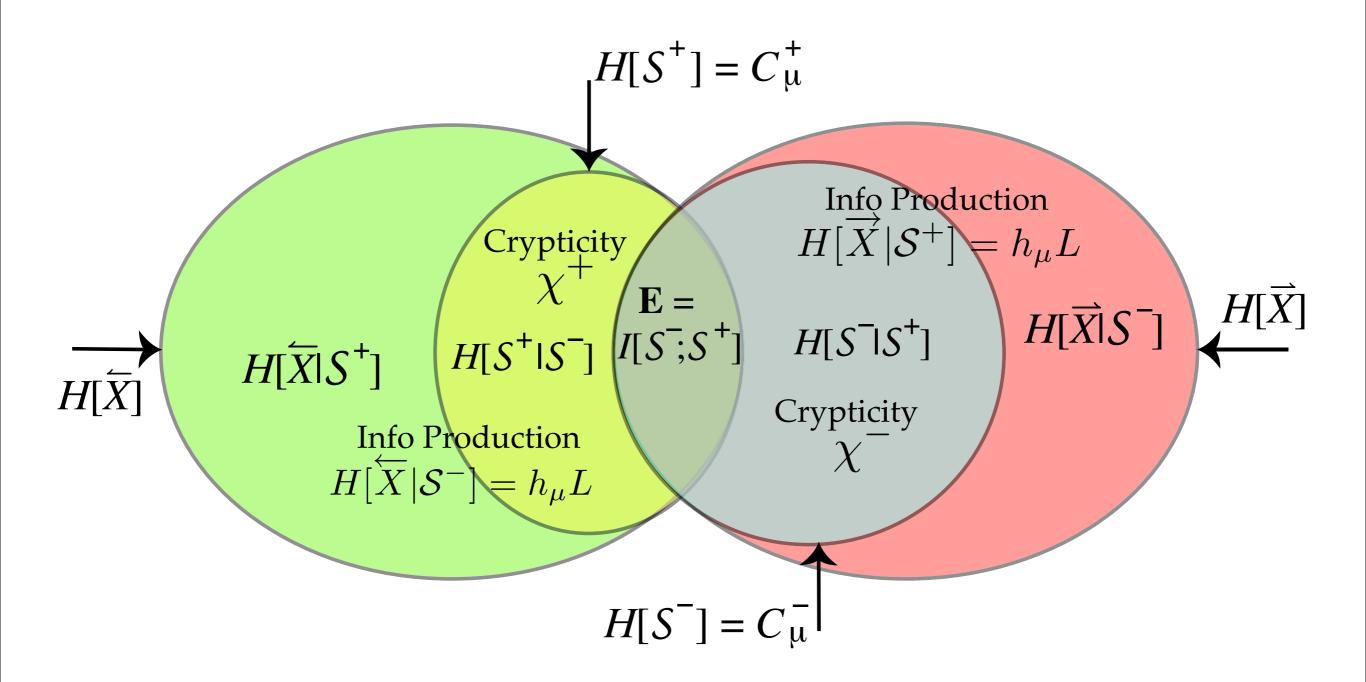
J. P. Crutchfield, C. J. Ellison, and J. R. Mahoney, "Time's Barbed Arrow: Irreversibility, Crypticity, and Stored Information", Physical Review Letters 103:9 (2009) 094101.











Labs Tonight

http://csc.ucdavis.edu/~cmg/