ϵM Reconstruction

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 $\epsilon \mathrm{M}$ Reconstruction:

Any method to go from process $\mathcal{P} \sim \Pr(\stackrel{\leftrightarrow}{S})$ to its ϵM

- (I) Analytical: Given model, equations of motion, description, ...
- (2) Statistical inference: Given samples of ${\mathcal P}$
 - (i) Subtree Reconstruction: Time or spacetime data to ϵM
 - (ii) State-splitting (CSSR):Time or spacetime data to ϵM
 - (iii) Spectral (eMSR): Power spectra to ϵM
 - (iv) Optimal Causal Inference: Time or spacetime data to ϵM

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How to reconstruct an ϵM : Subtree algorithm

Given: Word distributions $\Pr(s^D)$, $D=1,2,3,\ldots$

Steps:

- (I) Form depth-D parse tree.
- (2) Calculate node-to-node transition probabilities.
- (3) Causal states: Find morphs $\Pr(\overrightarrow{s}^L \mid \overrightarrow{s}^K)$ as subtrees.
- (4) Label tree nodes with morph (causal state) names.
- (5) Extract state-to-state transitions from parse tree.

(6) Assemble into
$$\epsilon M : \mathcal{M} = \left\{ \mathcal{S}, \{T^{(s)}, s \in \mathcal{A}\} \right\}$$
.

Algorithm parameters: D, L, K

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How to reconstruct an $\epsilon M...$

Examples:

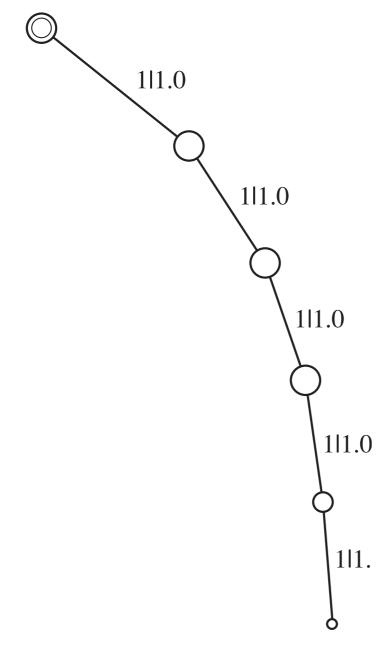
- (I) Period-I
- (2) Fair Coin
- (3) Period-2
- (4) Golden Mean Process
- (5) Even Process

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Examples (back to the Prediction Game):

Period-I: ...111111111111111

Parse Tree D = 5



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Examples (back to the Prediction Game):

Period-I: ...111111111111111

Parse Tree
$$D = 5$$

Morph
$$L = 2$$

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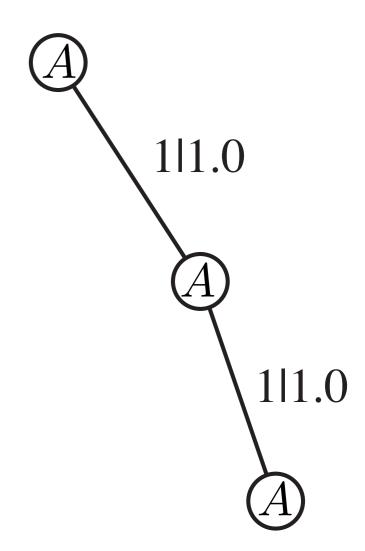
Examples (back to the Prediction Game) ...

Period-I: ...111111111111111

Space of histories: A single point.

One future morph: $\{1^+\}$

Morph A distribution: $\Pr(\stackrel{\rightarrow}{S}^L \mid \stackrel{\leftarrow}{s}) = 1$



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Examples (back to the Prediction Game) ...

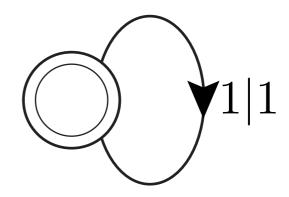
Period-I ...

$$\epsilon \mathbf{M} \colon \mathcal{M} = \left\{ \mathbf{S}, \left\{ T^{(s)}, s \in \mathcal{A} \right\} \right\}$$

$$S = \{S_0 = ...1111111\}$$

$$T^{(0)} = (0)$$

$$T^{(1)} = (1)$$

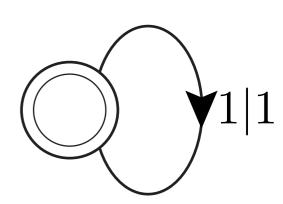


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Examples (back to the Prediction Game) ...

Period-I ...

Causal state distribution: $p_S = (1)$



Entropy Rate: $h_{\mu} = 0$ bits per symbol

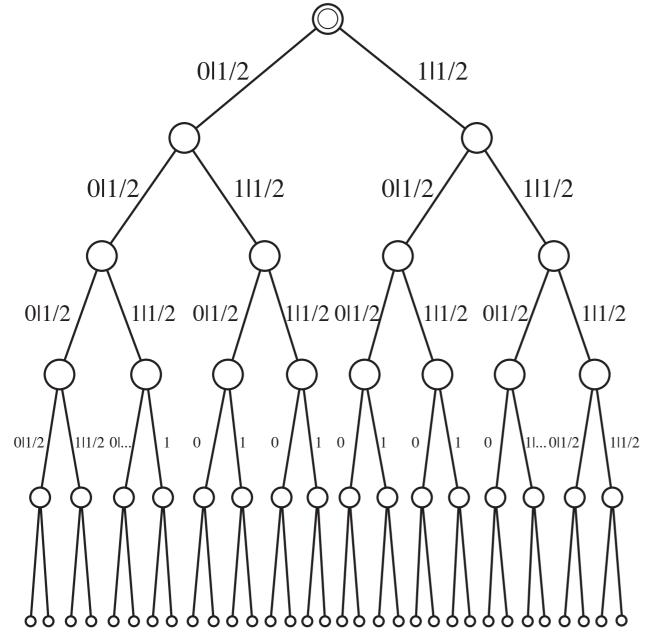
Statistical Complexity: $C_{\mu} = 0$ bits

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Examples (back to the Prediction Game) ...

Fair Coin: . . . 0101001110001101

Parse Tree D = 5



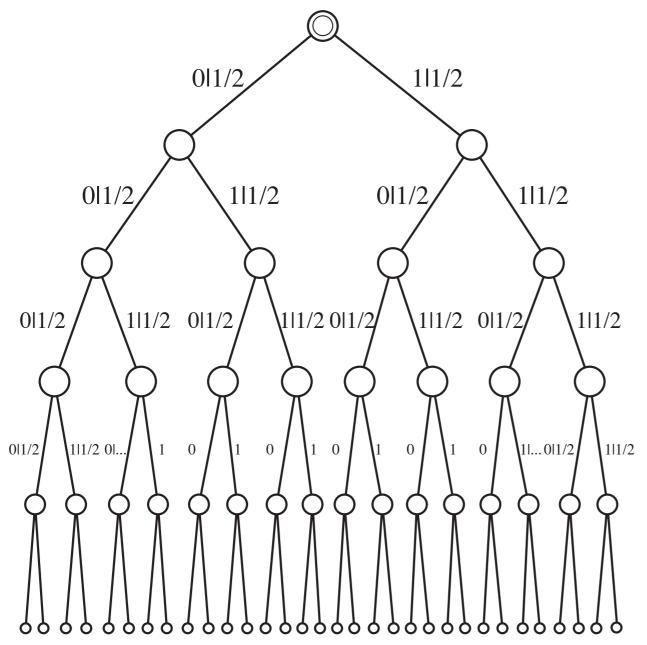
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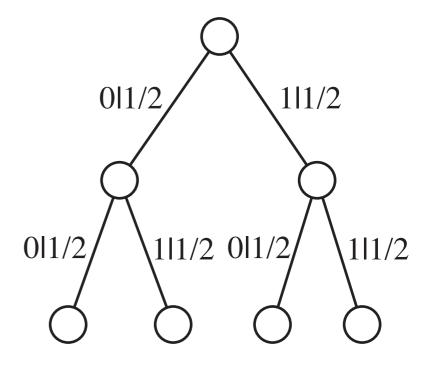
Examples (back to the Prediction Game) ...

Fair Coin: . . . 0101001110001101

Parse Tree D = 5

Morph L = 2





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Examples (back to the Prediction Game) ...

Fair Coin ...

Space of histories: $\overset{\leftarrow}{\mathbf{S}}^K = \mathcal{A}^K$

One future morph: \mathcal{A}^L

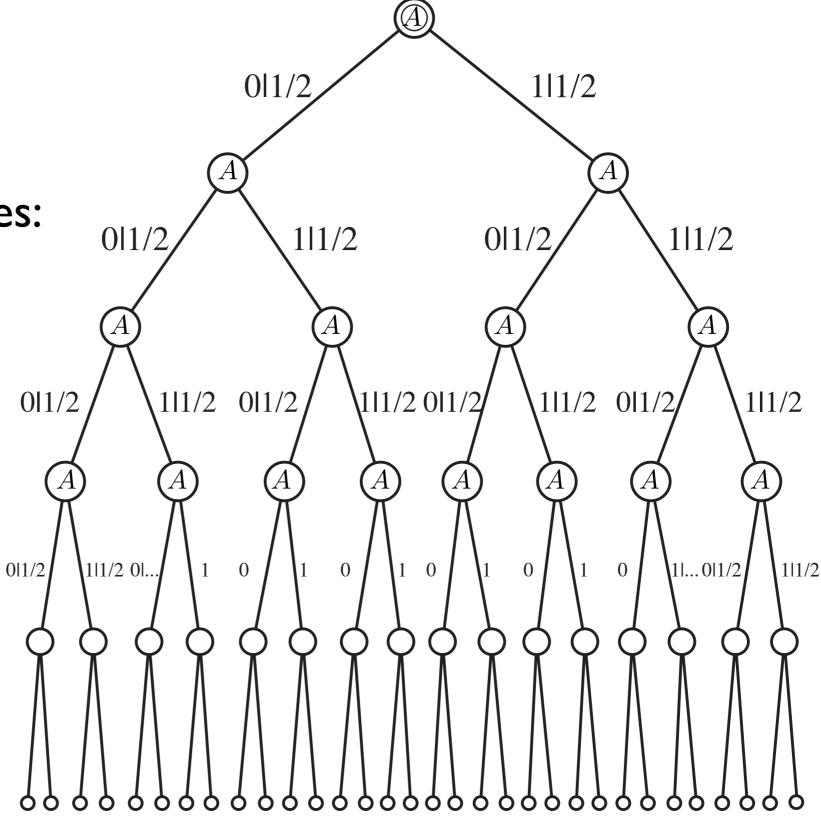
Morph A distribution: $\Pr(\vec{S}^L | \overset{\leftarrow}{s}) = 2^{-L}$

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Examples (back to the Prediction Game) ...

Fair Coin ...

Label tree nodes with state names:



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Examples (back to the Prediction Game) ...

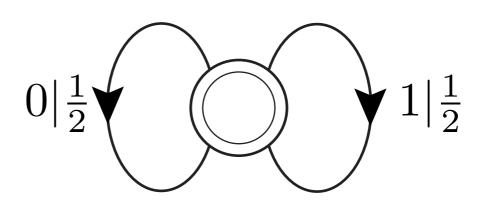
Fair Coin ...

$$\epsilon \mathbf{M} : \mathcal{M} = \left\{ \mathbf{S}, \left\{ T^{(s)}, s \in \mathcal{A} \right\} \right\}$$

$$\mathcal{S} = \{\mathcal{S}_0 = \mathcal{A}^L\}$$

$$T^{(0)} = \left(\frac{1}{2}\right)$$

$$T^{(1)} = \left(\frac{1}{2}\right)$$



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Examples (back to the Prediction Game) ...

Fair Coin ...

Causal state distribution: $p_S = (1)$

Entropy Rate: $h_{\mu} = 1$ bit per symbol

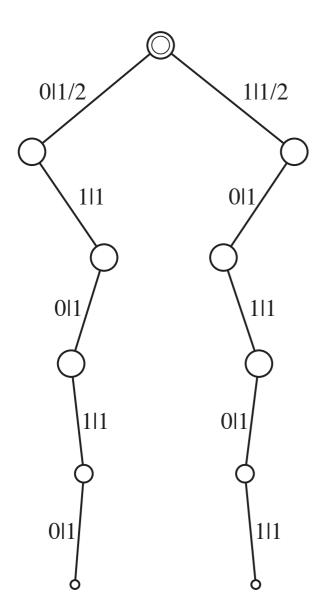
Statistical Complexity: $C_{\mu} = 0$ bits

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Examples (back to the Prediction Game) ...

Period-2 Process: . . . 01010101010101

Parse Tree D = 5



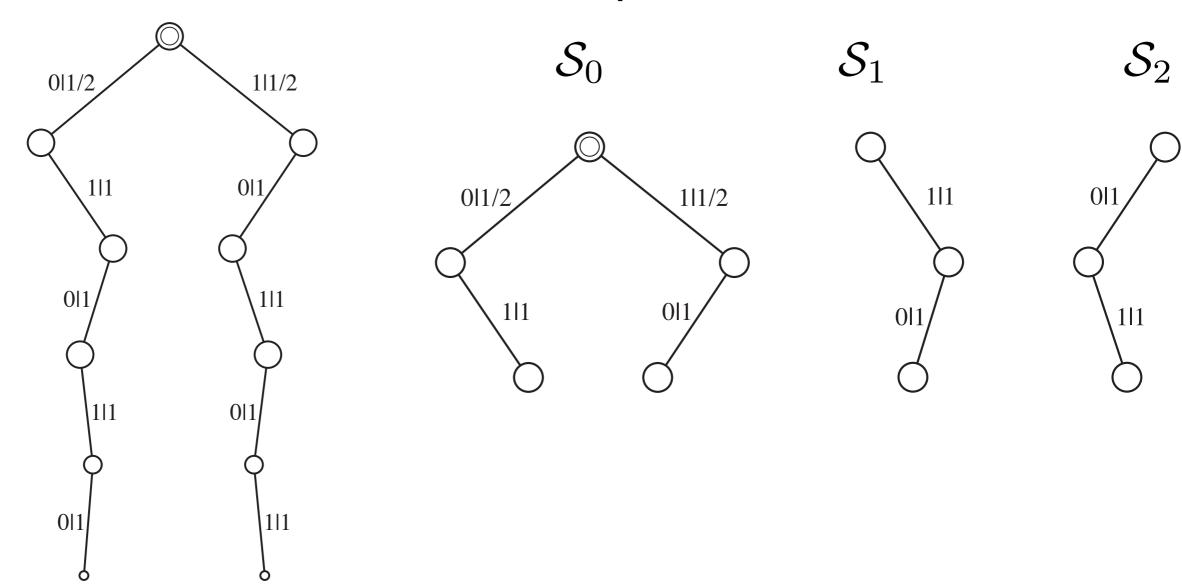
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Examples (back to the Prediction Game) ...

Period-2 Process: . . . 01010101010101

Parse Tree D = 5

Morphs at L = 2



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Examples (back to the Prediction Game) ...

Morph distributions:

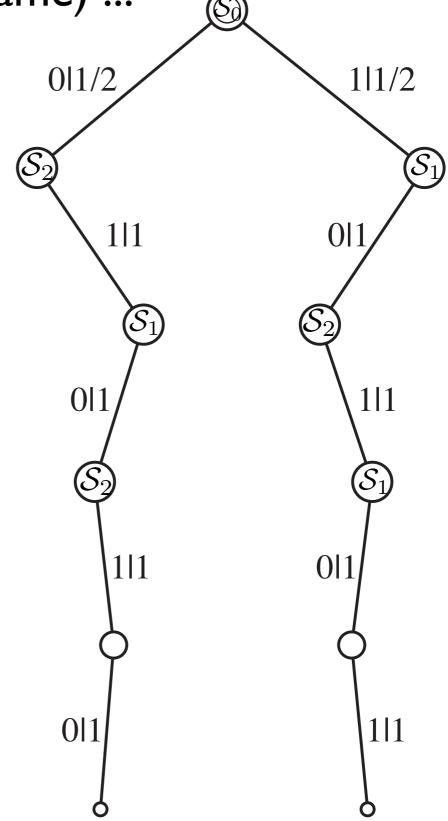
$$\Pr(0|\lambda) = \frac{1}{2}$$
 $\Pr(1|0) = 1$ $\Pr(0|0) = 0$ $\Pr(1|\lambda) = \frac{1}{2}$ $\Pr(1|1) = 0$ $\Pr(0|1) = 1$

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Examples (back to the Prediction Game) ...

Period-2 Process ...

Label tree nodes:



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Examples (back to the Prediction Game) ...

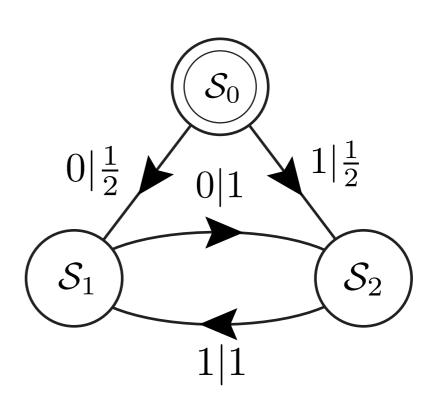
Period-2 Process ...

$$\mathcal{M} = \left\{ \mathcal{S}, \left\{ T^{(s)}, s \in \mathcal{A} \right\} \right\}$$

$$S = \{S_0 = \{...0101, ...1010\}, S_1 = \{...1010\}, S_2 = \{...0101\}\}$$

$$T^{(0)} = \begin{pmatrix} 0 & \frac{1}{2} & 0 \\ 0 & 0 & 1 \\ 0 & 0 & 0 \end{pmatrix}$$

$$T^{(1)} = \begin{pmatrix} 0 & 0 & \frac{1}{2} \\ 0 & 0 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$



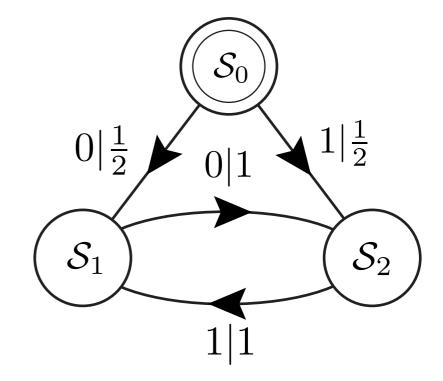
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Examples (back to the Prediction Game) ...

Period-2 Process ...

Causal State Distribution:

$$p_{\mathcal{S}} = \left(0, \frac{1}{2}, \frac{1}{2}\right)$$



Entropy rate: $h_{\mu} = 0$ bits per symbol

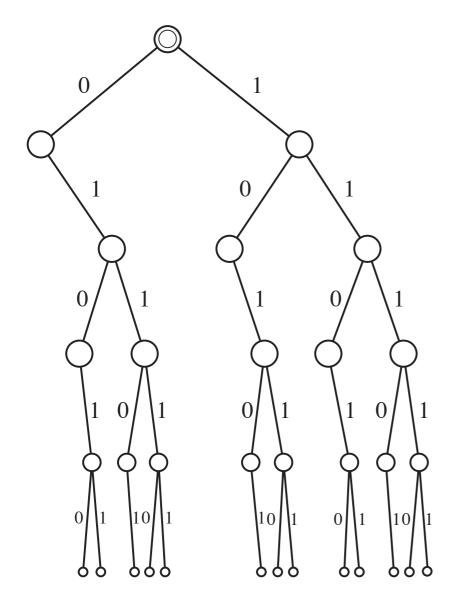
Statistical complexity: $C_{\mu} = 1 \ \mathrm{bit}$

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Examples ...

Golden Mean Process:

Parse Tree D = 5



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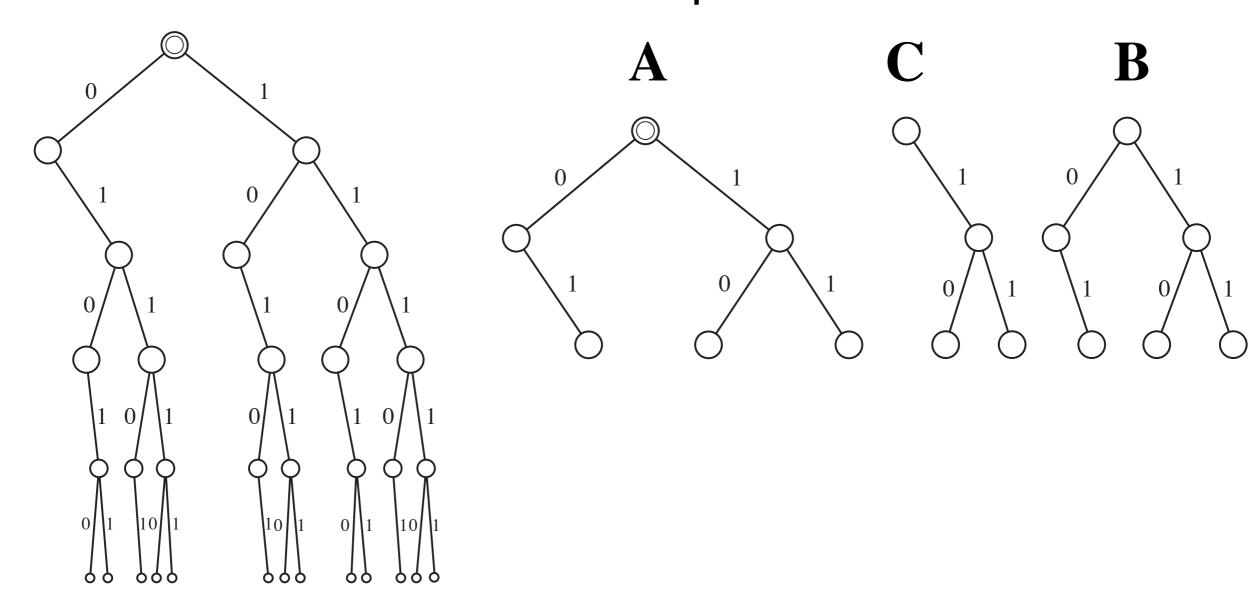
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Examples ...

Golden Mean Process: Topological Reconstruction

Parse Tree D = 5

Morphs at L = 2



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Examples ...

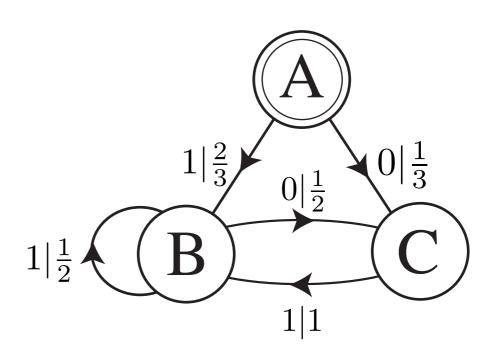
Golden Mean Process ...

$$\mathcal{M} = \left\{ \mathcal{S}, \left\{ T^{(s)}, s \in \mathcal{A} \right\} \right\}$$

$$\mathcal{S} = \{A, B, C\}$$

$$T^{(0)} = \begin{pmatrix} 0 & 0 & 1/3 \\ 0 & 0 & 1/2 \\ 0 & 0 & 0 \end{pmatrix}$$

$$T^{(1)} = \begin{pmatrix} 0 & 2/3 & 0 \\ 0 & 1/2 & 0 \\ 0 & 1 & 0 \end{pmatrix}$$



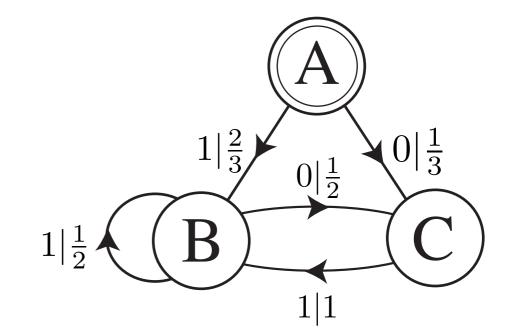
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Examples ...

Golden Mean Process ...

Causal State Distribution:

$$p_{\mathcal{S}} = \left(0, \frac{2}{3}, \frac{1}{3}\right)$$



Entropy rate: $h_{\mu} = \frac{2}{3}$ bits per symbol

Statistical complexity: $C_{\mu} = H(\frac{2}{3})$ bits

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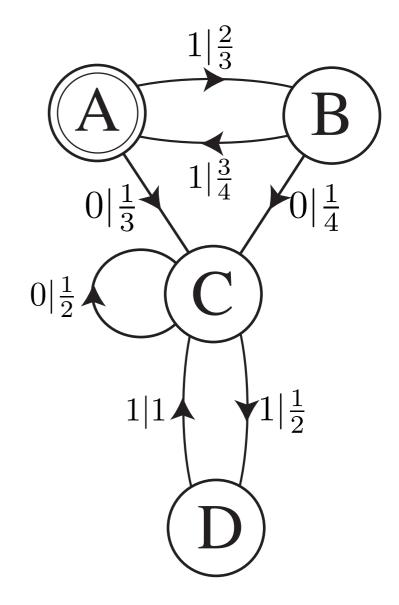
Examples ...

Even Process:

$$\mathcal{M} = \left\{ \mathcal{S}, \left\{ T^{(s)}, s \in \mathcal{A} \right\} \right\}$$

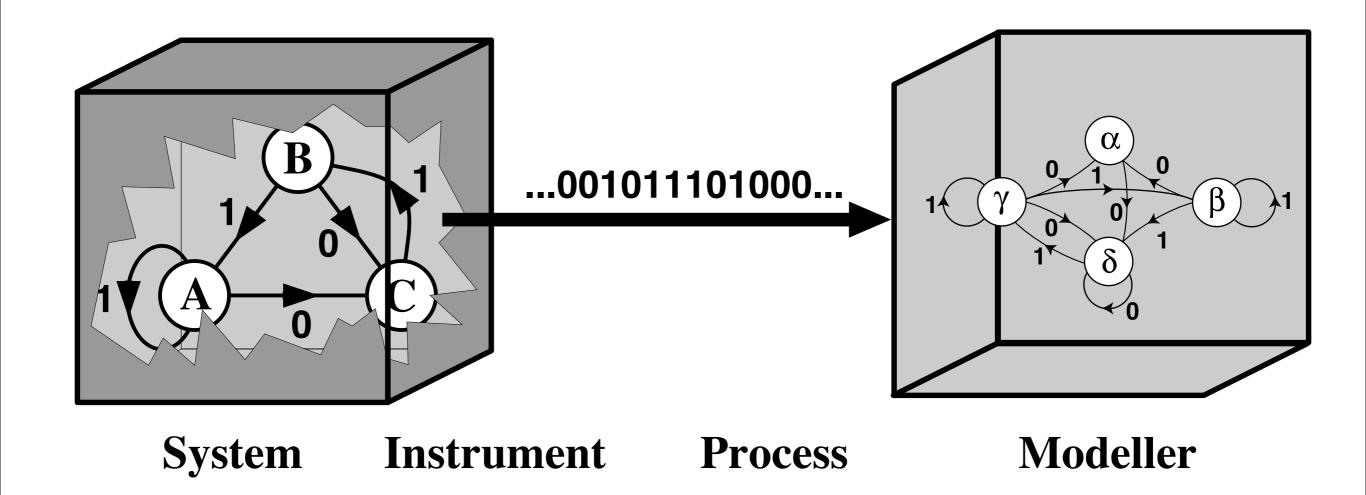
Entropy rate: $h_{\mu} = \frac{2}{3}$ bits per symbol

Statistical complexity: $C_{\mu} = H(\frac{2}{3})$ bits



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The Learning Channel:



Central questions:

What are the states? Causal States

What is the dynamic? The ϵ -Machine

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Process \Rightarrow Predictive equivalence \Rightarrow ϵ – Machine

$$\Pr(\overset{\leftrightarrow}{S}) \Rightarrow \overset{\leftarrow}{\mathbf{S}} / \sim \Rightarrow \epsilon - \text{Machine}$$

$$\mathcal{M} = \left\{ \mathcal{S}, \left\{ T^{(s)}, s \in \mathcal{A} \right\} \right\}$$

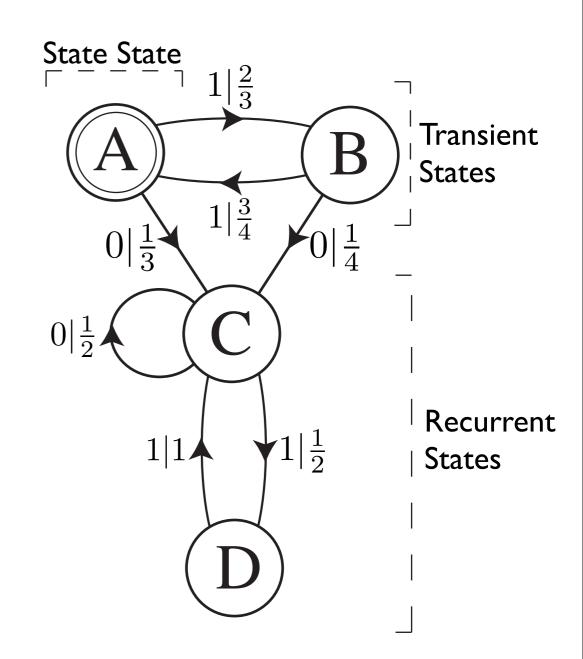
Unique Start State:

$$\mathcal{S}_0 = [\lambda]$$

 $\Pr(\mathcal{S}_0, \mathcal{S}_1, \mathcal{S}_2, \ldots) = (1, 0, 0, \ldots)$

Transient States

Recurrent States



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A Model of a Process $\Pr(\stackrel{\leftrightarrow}{S})$:

ϵ -Machine reproduces the process's word distribution:

$$Pr(s^{1}), Pr(s^{2}), Pr(s^{3}), \dots$$

$$s^{L} = s_{1}s_{2}\dots s_{L} \qquad \mathcal{S}(t=0) = \mathcal{S}_{0}$$

$$Pr(s^{L}) = Pr(\mathcal{S}_{0})Pr(\mathcal{S}_{0} \to_{s=s_{1}} \mathcal{S}(1))Pr(\mathcal{S}(1) \to_{s=s_{2}} \mathcal{S}(2))$$

$$\dots Pr(\mathcal{S}(L-1) \to_{s=s_{L}} \mathcal{S}(L))$$

Initially, $\Pr(\mathcal{S}_0) = 1$.

$$\Pr(s^{L}) = \prod_{l=1}^{L} T_{i=\epsilon(s^{l-1}), j=\epsilon(s^{l})}^{(s_{l})}$$

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A Model of a Process $\Pr(\stackrel{\longleftrightarrow}{S})$...

Calculate word distribution from recurrent states: $S_i \in S_{ ext{recurrent}}$

$$\Pr(s^1), \Pr(s^2), \Pr(s^3), \dots$$

Get
$$\langle \pi | = (p_{\mathcal{S}_1}, p_{\mathcal{S}_2}, \ldots)$$
 from $T = \sum_{s \in A} T^{(s)}$

Then

$$\Pr(s) = \langle \pi | T^{(s)} | 1 \rangle \qquad |1\rangle = \begin{pmatrix} 1 \\ 1 \\ \vdots \end{pmatrix}$$

$$\Pr(s_0 s_1) = \langle \pi | T^{(s_0)} T^{(s_1)} | 1 \rangle$$

• • •

$$\Pr(s^L) = \langle \pi | T^{(s^L)} | 1 \rangle$$

$$T^{(s^L)} = T^{(s_0)}T^{(s_1)}\cdots T^{(s_{L-1})}$$

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Properties:

Conditional independence of future & past

Unifilar

Markovian

Optimal predictor

Minimal size

Unique

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Causal shielding:

Past and future are independent given causal state

Process:
$$\Pr(\stackrel{\leftrightarrow}{S}) = \Pr(\stackrel{\leftarrow}{S}\stackrel{\rightarrow}{S})$$

$$\Pr(\stackrel{\leftarrow}{S}\stackrel{\rightarrow}{S}|\mathcal{S}) = \Pr(\stackrel{\leftarrow}{S}|\mathcal{S})\Pr(\stackrel{\rightarrow}{S}|\mathcal{S})$$

Causal states shield past & future from each other.

Similar to states of a Markov chain, but for hidden processes.

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Proof sketch:

$$\Pr(\stackrel{\leftrightarrow}{S} | \mathcal{S}) = \Pr(\stackrel{\leftarrow}{S} \stackrel{\rightarrow}{S} | \mathcal{S})$$

$$= \Pr(\stackrel{\rightarrow}{S} | \stackrel{\leftarrow}{S}, \mathcal{S}) \Pr(\stackrel{\leftarrow}{S} | \mathcal{S})$$

Will show: $\Pr(\overrightarrow{S} \mid \overleftarrow{S}, \mathcal{S}) = \Pr(\overrightarrow{S} \mid \mathcal{S})$

$$S = \epsilon(\overleftarrow{s}) \Rightarrow$$

$$\Pr\left(\overrightarrow{S} \middle| \overleftarrow{S} = \overleftarrow{s}', S = \epsilon(\overleftarrow{s})\right) = \Pr\left(\overrightarrow{S} \middle| \overleftarrow{S} = \overleftarrow{s}\right) \quad (\overleftarrow{s}' \in [\overleftarrow{s}])$$

But, also,
$$\Pr\left(\overrightarrow{S}\mid \overrightarrow{S}=\overleftarrow{s}\right)=\Pr\left(\overrightarrow{S}\mid \mathcal{S}=\epsilon(\overleftarrow{s})\right)$$
 (Causal equiv. rel'n)

So,
$$\Pr\left(\overrightarrow{S} \mid \overleftarrow{S} = \overleftarrow{s}, \mathcal{S} = \sigma\right) = \Pr\left(\overrightarrow{S} \mid \mathcal{S} = \sigma\right)$$

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 ϵMs are Unifilar: $(\mathcal{S}_t, s) \rightarrow \text{unique } \mathcal{S}_{t+1}$

(in automata theory, "deterministic")

That is:

(I) $S_i \in \mathcal{S}, \ s \in \mathcal{A}$, at most one $S_j \in \mathcal{S}$:

$$\stackrel{\leftarrow}{s} \in \mathcal{S}_i \Rightarrow \stackrel{\leftarrow}{s} s \in \mathcal{S}_j$$

(2) If there is a next causal state j:

$$S_{k \neq j} \in \mathcal{S} \Rightarrow T_{ik}^{(s)} = 0$$

(3) If there is not:

$$T_{ij}^{(s)} = 0$$

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Unifilarity ...

Proof sketch:

Must show $\overset{\leftarrow}{s} \sim \overset{\leftarrow}{s}' \Rightarrow \overset{\leftarrow}{s}s \sim \overset{\leftarrow}{s}'s$

Futures with symbol prefixed: sF

$$F \subset \mathcal{A}^{\infty}$$

$$\begin{array}{ccc}
\overleftarrow{S} \sim \overleftarrow{S}' & \Rightarrow & \Pr\left(\overrightarrow{S} \in sF | \overleftarrow{S} = \overleftarrow{s}\right) = \Pr\left(\overrightarrow{S} \in sF | \overleftarrow{S} = \overleftarrow{s}'\right) \\
& \Pr\left(\overrightarrow{S}^1 = s, \overrightarrow{S}_1 \in F | \overleftarrow{S} = \overleftarrow{s}\right) = \Pr\left(\overrightarrow{S}^1 = s, \overrightarrow{S}_1 \in F | \overleftarrow{S} = \overleftarrow{s}'\right) \\
& \Pr\left(\overrightarrow{S}_1 \in F | \overrightarrow{S}^1 = s, \overleftarrow{S} = \overleftarrow{s}\right) \Pr\left(\overrightarrow{S}^1 = s | \overleftarrow{S} = \overleftarrow{s}'\right) = \Pr\left(\overrightarrow{S}_1 \in F | \overrightarrow{S}^1 = s, \overleftarrow{S} = \overleftarrow{s}'\right) \Pr\left(\overrightarrow{S}^1 = s | \overleftarrow{S} = \overleftarrow{s}'\right) \\
& \Pr\left(\overrightarrow{S}_1 \in F | \overleftarrow{S} = \overleftarrow{s}'\right) = \Pr\left(\overrightarrow{S}_1 \in F | \overleftarrow{S} = \overleftarrow{s}'s\right) & \text{(Stationarity and by assumption} \\
& \Rightarrow \overleftarrow{s} s \sim \overleftarrow{s}' s & \bigcirc & \Pr(\overrightarrow{S}^1 = s | \overleftarrow{s}) = 1 \\
& \Pr(\overrightarrow{S}^1 = s | \overleftarrow{s}') = 1
\end{array}\right)$$

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The ϵ -Machine ... Unifiliarity ...

Consequence:

Unifilarity: I-I map between state-sequences & symbol-sequences.

Entropy rate expression requires this I-I mapping.

Can (must) use ϵM to calculate entropy rate h_{μ} .

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$\epsilon \mathrm{Ms}$ are Optimal Predictors:

Compared to any rival effective states R:

$$H\left[\overrightarrow{S}^{L}|R\right] \ge H\left[\overrightarrow{S}^{L}|\mathcal{S}\right]$$

Proof sketch:
$$H\begin{bmatrix} \overrightarrow{S}^L | \mathcal{S} \end{bmatrix} = H\begin{bmatrix} \overrightarrow{S}^L | \overleftarrow{s} \in \mathcal{S} \end{bmatrix}$$
 (Causal equiv. rel'n)
$$= H\begin{bmatrix} \overrightarrow{S}^L | \overleftarrow{s} \end{bmatrix}$$

$$\leq H\begin{bmatrix} \overrightarrow{S}^L | R \end{bmatrix} \qquad R = \eta(\overleftarrow{s})$$
 (Data processing inequality)

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 $\epsilon \mathrm{Ms}$ are Optimal Predictors ...

Lemma:

$$h_{\mu}(\mathcal{S}) = h_{\mu}$$

Proof:
$$h_{\mu}(S) = \lim_{L \to \infty} \frac{1}{L} H \left[\overrightarrow{S}^{L} | \mathcal{S} \right]$$

$$= \lim_{L \to \infty} \frac{1}{L} H \left[\overrightarrow{S}^{L} | \overleftarrow{S} \right]$$

$$= \lim_{L \to \infty} \frac{1}{L} L H \left[S | \overleftarrow{S} \right]$$

$$= H \left[S | \overleftarrow{S} \right]$$

$$= h_{\mu}$$

(Block entropy)

(Causal equiv. rel'n)

(Stationarity)

Corollary: $h_{\mu}(R) \geq h_{\mu}$

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 $\epsilon \mathrm{Ms}$ are Optimal Predictors ...

Corollary (Maximal Prescience): $\Pi(R) \leq \Pi(S)$

Rival model: $\Pi(R) = \log_2 |\mathcal{A}| - h_{\mu}(R)$

But: $\Pi(S) = \log_2 |A| - h_\mu = G$

So: $\Pi(R) \leq \Pi(S)$

 $h_{\mu}(R) \geq h_{\mu}$

Remarks:

- (I) Causal states contain every difference (in past) that makes a difference (to future) (Bateson "information")
- (2) Causal states are sufficient statistics for the future. (See below.)

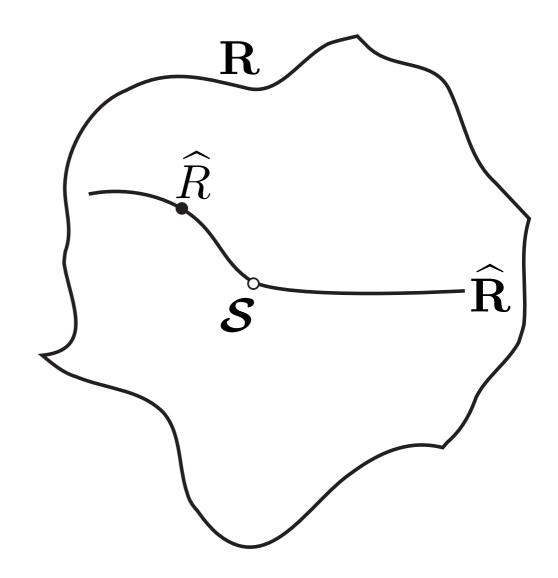
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Prescient Rivals $\widehat{\mathbf{R}}$:

Alternative models that are optimal predictors

$$\widehat{R} \in \widehat{\mathbf{R}}$$

$$H[\overrightarrow{S}^{L} | \widehat{R}] = H[\overrightarrow{S}^{L} | \mathcal{S}]$$



(Prescient rivals are sufficient statistics for process's future.)

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Minimal Statistical Complexity:

For all prescient rivals, ϵM is the smallest:

$$C_{\mu}(\widehat{R}) \ge C_{\mu}(\mathcal{S})$$

Proof sketch:

(I) Prescient rivals are refinements, so

$$\exists g: \mathcal{S} = g(\widehat{R})$$

(2) But

$$H[f(X)] \le H[X] \Rightarrow H[S] = H[g(\widehat{R})] \le H[\widehat{R}]$$

(3) So
$$C_{\mu} \leq H[\widehat{R}]$$



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Minimal Statistical Complexity ...

Consequence:

- (I) C_{μ} measures historical information process stores.
- (2) This would not be true, if not minimal representation.

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Summary:

 ϵM :

- (I) Optimal predictor: Lower prediction error than any rival.
- (2) Minimal size: Smallest of the prescient rivals.
- (3) Unique: Smallest, optimal, unifilar predictor is equivalent.
- (4) Model of the process: Reproduces all of process's statistics.
- (5) Causal shielding: Renders process's future independent of past.

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Dynamical system's intrinsic computation:

- (I) How much of past does process store?
- (2) In what architecture is that information stored?
- (3) How is stored information used to produce future behavior?

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