Information Theory: Part II Applications to Stochastic Processes

 We now consider applying information theory to a long sequence of measurements.

 $\cdots 00110010010101101001100111010110 \cdots$

- In so doing, we will be led to two important quantities
 - 1. Entropy Rate: The irreducible randomness of the system.
 - 2. Excess Entropy: A measure of the complexity of the sequence.

Context: Consider a long sequence of discrete random variables. These could be:

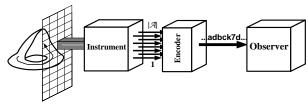
- 1. A long time series of measurements
- 2. A symbolic dynamical system
- 3. A one-dimensional statistical mechanical system

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The Measurement Channel

 Can also picture this long sequence of symbols as resulting from a generalized measurement process:



- On the left is "nature"—some system's state space.
- The act of measurement projects the states down to a lower dimension and discretizes them.
- The measurements may then be encoded (or corrupted by noise).
- They then reach the observer on the right.
- Figure source: Crutchfield, "Knowledge and Meaning ... Chaos and Complexity." In Modeling Complex Systems. L. Lam and H. C. Morris, eds. Springer-Verlag, 1992: 66-10.

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Stochastic Process Notation

- Random variables $S_i, S_i = s \in \mathcal{A}$.
- Infinite sequence of random variables: $\overset{\leftrightarrow}{S} = \dots S_{-1} S_0 S_1 S_2 \dots$
- Block of L consecutive variables: $S^L = S_1, \ldots, S_L$.
- $\Pr(s_i, s_{i+1}, \dots, s_{i+L-1}) = \Pr(s^L)$
- Assume translation invariance or stationarity:

$$\Pr(s_{i}, s_{i+1}, \dots, s_{i+L-1}) = \Pr(s_{1}, s_{2}, \dots, s_{L}).$$

- Left half ("past"): $\overset{\leftarrow}{S} \equiv \cdots S_{-3} \, S_{-2} \, S_{-1}$
- ullet Right half ("future"): $\overset{
 ightarrow}{S} \equiv S_0 \, S_1 \, S_2 \cdots$

 $\cdots 11010100101101010101001001010010 \cdots$

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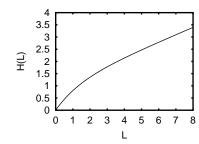
Entropy Growth

• Entropy of *L*-block:

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$$H(L) \equiv -\sum_{s^L \in \mathcal{A}^L} \Pr(s^L) \log_2 \Pr(s^L) .$$

ullet H(L)= average uncertainty about the outcome of L consecutive variables.



- ullet H(L) increases monotonically and asymptotes to a line
- We can learn a lot from the shape of H(L).

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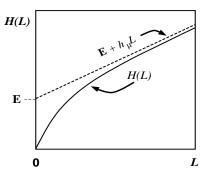
• $h_{\mu}(L) = H[S_L|S_1S_1...S_{L-1}]$

previous L symbols have been observed.

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Entropy Rate

• Let's first look at the slope of the line:



- Slope of H(L): $h_{\mu}(L) \equiv H(L) H(L-1)$
- Slope of the line to which H(L) asymptotes is known as the *entropy rate*:

$$h_{\mu} = \lim_{L \to \infty} h_{\mu}(L).$$

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Interpretations of Entropy Rate

- Uncertainty per symbol.
- Irreducible randomness: the randomness that persists even after accounting for correlations over arbitrarily large blocks of variables.
- The randomness that cannot be "explained away".
- Entropy rate is also known as the Entropy Density or the Metric Entropy.
- $h_{\mu} =$ Lyapunov exponent for many classes of 1D maps.
- ullet The entropy rate may also be written: $h_{\mu}=\lim_{L
 ightarrow \infty} rac{H(L)}{L}$.
- ullet h_{μ} is equivalent to thermodynamic entropy.
- These limits exist for all stationary processes.

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How does $h_{\mu}(L)$ approach h_{μ} ?

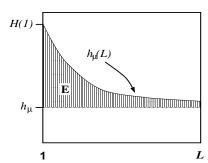
Entropy Rate, continued

• Slope of the line to which H(L) asymptotes is known as the *entropy rate*:

ullet I.e., $h_{\mu}(L)$ is the average uncertainty of the next symbol, given that the

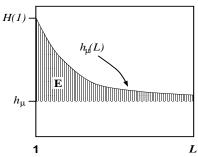
 $h_{\mu} = \lim_{L \to \infty} h_{\mu}(L).$

ullet For finite L , $h_{\mu}(L) \geq h_{\mu}$. Thus, the system appears more random than it is.



• We can learn about the complexity of the system by looking at *how* the entropy density converges to h_{μ} .

The Excess Entropy



• The excess entropy captures the nature of the convergence and is defined as the shaded area above:

$$\mathbf{E} \equiv \sum_{L=1}^{\infty} [h_{\mu}(L) - h_{\mu}] .$$

• E is thus the total amount of randomness that is "explained away" by considering larger blocks of variables.

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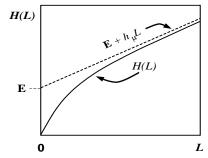
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Excess Entropy: Other expressions and interpretations

Geometric View

- **E** is the *y*-intercept of the straight line to which H(L) asymptotes.
- $\mathbf{E} = \lim_{L \to \infty} [H(L) h_{\mu}L]$.



Excess Entropy: Other expressions and interpretations

Mutual information

• One can show that **E** is equal to the mutual information between the "past" and the "future":

$$\mathbf{E} = I(\overset{\leftarrow}{S}; \vec{S}) \equiv \sum_{\overset{\leftarrow}{\{\vec{s}\}}} \Pr(\vec{s}) \log_2 \left[\frac{\Pr(\vec{s})}{\Pr(\overset{\leftarrow}{s}) \Pr(\vec{s})} \right] .$$

- E is thus the amount one half "remembers" about the other, the reduction in uncertainty about the future given knowledge of the past.
- Equivalently, E is the "cost of amnesia:" how much more random the future appears if all historical information is suddenly lost.

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Excess Entropy Summary

- Is a structural property of the system measures a feature complementary to entropy.
- Measures memory or spatial structure.
- Lower bound for statistical complexity, minimum amount of information needed for minimal stochastic model of system

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2.5

0.5

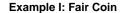
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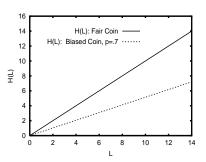
• Sequence: ... 1010111011101110...

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15

16





- For fair coin, $h_{\mu}=1$.
- For the biased coin, $h_{\mu} \approx 0.8831$.
- For both coins, $\mathbf{E} = 0$.
- Note that two systems with different entropy rates have the same excess entropy.

David P. Feldman

http://hornacek.coa.edu/dave

David P. Feldman

http://hornacek.coa.edu/dave

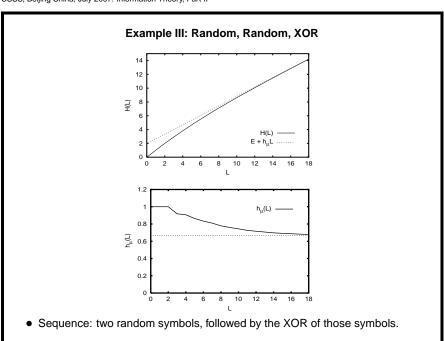
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Example II, continued

- Sequence: ... 1010111011101110...
- $h_{\mu} \approx 0$; the sequence is perfectly predictable.
- ullet $\mathbf{E} = \log_2 16 = 4$: four bits of phase information
- For any period-p sequence, $h_{\mu} = 0$ and $\mathbf{E} = \log_2 p$.

For more than you probably ever wanted to know about periodic sequences, see Feldman and Crutchfield, Synchronizing to Periodicity: The Transient Information and Synchronization Time of Periodic Sequences. Advances in Complex Systems. 7(3-4): 329-355, 2004.

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Example II: Periodic Sequence

10 12 14 16

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Example III. continued

- Sequence: two random symbols, followed by the XOR of those symbols.
- $h_{\mu} = \frac{2}{3}$; two-thirds of the symbols are unpredictable.
- $\mathbf{E} = \log_2 4 = 2$: two bits of phase information.
- For many more examples, see Crutchfield and Feldman, Chaos, 15: 25-54, 2003.

David P. Feldman

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Excess Entropy: Selected References and Applications

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- Szépfalusy and Györgyi, *Phys. Rev. A*, 33:2852-2855, 1986. [Dynamical systems]
- Lindgren and Nordahl, Complex Systems, 2:409-440. (1988). [CAs, dynamical
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- Feldman and Crutchfield, SFI:98-04-026, 1998. Crutchfield and Feldman, Phys. Rev. E 55:R1239-42. 1997. [One-dimensional Ising models]

Excess Entropy: Notes on Terminology

All of the following terms refer to essentially the same quantity.

• Excess Entropy: Crutchfield, Packard, Feldman

• Stored Information: Shaw

• Effective Measure Complexity: Grassberger, Lindgren, Nordahl

• Reduced (Rényi) Information: Szépfalusy, Györgyi, Csordás

• Complexity: Li, Arnold

• Predictive Information: Nemenman, Bialek, Tishby

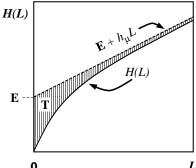
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- Ebeling. Physica D, 1090:42-52. 1997. [Dynamical systems, written texts, music]
- Bialek, et al, Neur. Comp., 13:2409-2463. 2001. [Long-range 1D Ising models, machine learning]

Transient Information T

- $\mathbf{T} \equiv \sum_{L=1}^{\infty} [\mathbf{E} + h_{\mu}L H(L)].$
- T is related to the total uncertainty experienced while synchronizing to a process.



- The shaded area is the transient information T.
- T measures how difficult it is to synchronize to a sequence.

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Some Applications in Agent-Based Modeling Settings, continued

- 3. More generally it seems likely that the entropy and mutual information are useful tools for quantifying
 - (a) properties of agents: e.g., how much memory they have
 - (b) the behavior of agents: e.g, how unpredictably they act
 - (c) properties of the environment: e.g., how structured it is

Some Applications in Agent-Based Modeling Settings

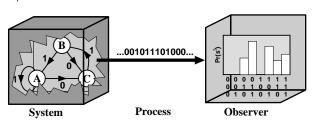
- 1. If an agent doesn't have sufficient memory, its environment will appear more random. In a quantitative sense, regularities that are missed (as measured by the excess entropy) are converted into randomness (as measured by the entropy rate).
 - Crutchfield and Feldman, Synchronizing to the Environment: Information Theoretic Constraints on Agent Learning. Advances in Complex Systems. 4. 251-264. 2001.
- 2. The average-case difficulty for an agent to synchronize to a periodic environment is measured by the transient information.
 - Feldman and Crutchfield. Synchronizing to a Periodic Signal: The Transient Information and Synchronization Time of Periodic Sequences. Advances in Complex Systems. 7. 329-355. 2004.

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Estimating Probabilities

ullet E and h_{μ} can be estimated empirically by observing a process.



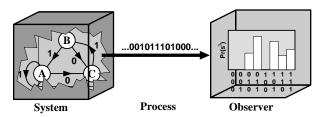
• One simply forms histograms of occurrences of particular sequences and uses these to estimate $\Pr(s^L)$, from which **E** and h_μ may be readily calculated.

For more sophisticated and accurate ways of inferring h_{μ} , see, e.g.,

- Schürmann and Grassberger. Chaos 6:414-427. 1996.
- Nemenman. http://arXiv.org/physics/0207009. 2002.

A look ahead

• Note that the observer sees measurement symbols: 0's and 1's.



- It doesn't see inside the "black box" of the system.
- $\bullet\,$ In particular, it doesn't see the internal, hidden states of the system, A,B, and C.
- Is there a way an observer can infer these hidden states?
- What is the meaning of state?

David P. Feldman

http://hornacek.coa.edu/dave